

```
int coin_collect(int a[][100], int n){  
    int i,j, coins[100][100];  
  
    coins[0][0] = a[0][0]; //initial cell  
  
    for (i=1; i<n; i++) //first row  
        coins[0][i] = coins[0][i-1] + a[0][i];  
  
    for (i=1; i<n; i++) //first column  
        coins[i][0] = coins[i-1][0] + a[i][0];  
  
    for (i=1; i<n; i++) //filling up the rest of the array  
        for (j=1; j<n; j++)  
            coins[i][j] = max(coins[i-1][j], coins[i][j-1])  
                + a[i][j];  
  
    return coins[n-1][n-1]; //value of last cell  
}
```

```
int max(int a, int b){  
    if (a>b) return a;  
    else return b;  
}  
  
int main(){  
    int m[100][100],i,j,n;  
  
    scanf("%d", &n);  
    for (i=0; i<n; i++)  
        for (j=0; j<n; j++)  
            scanf("%d", &m[i][j]);  
  
    printf("%d\n", coin_collect(m,n));  
    return 0;  
}
```

# Passing two dimensional arrays as parameters

Write a program that takes a two dimensional array of type double [5][6] and prints the sum of entries in each row.

```
void marginals(double mat[5][6]) {  
    int i,j; double rowsum;  
    for (i=0; i < 5; i=i+1) {  
        rowsum = 0.0;  
        for (j=0; j < 6; j = j+1) {  
            rowsum = rowsum+mat[i][j];  
        }  
        printf("%f ", rowsum);  
    }  
}
```

## Question?

But suppose we have read only the first 3 rows out of the 5 rows of mat. And we would like to find the marginal sum of the first 3 rows.

## Answer:

That's easy, we can take an additional parameter **nrows** and run the loop for  $i=0..(nrows-1)$  instead of  $0..5$ .

The slightly generalized program would be:

```
void marginals(double mat[5][6], int nrows) {  
    int i,j; double rowsum;  
    for (i=0; i < nrows; i=i+1) {  
        rowsum = 0.0;  
        for (j=0; j < 6; j = j+1) {  
            rowsum = rowsum+mat[i][j];  
        }  
        printf("%f ", rowsum);  
    }  
}
```

In parameter  
double mat[5][6],  
**C completely**  
ignores the  
number of rows 5.  
It is only  
interested in the  
number of cols: 6.

We declared mat to be of type double [5][6].  
Does this mean that nrows should be  $\leq 5$ ?  
We are not checking for it!

Let's see more  
examples...

The following program is exactly identical to the previous one.

```
void marginals(double mat[ ][6], int nrows)
{
    int i,j; int rowsum;
    for (i=0; i < nrows; i=i+1) {
        rowsum = 0.0;
        for (j=0; j < 6; j = j+1) {
            rowsum = rowsum+mat[i][j];
        }
        printf("%f ", rowsum);
    }
}
```

This means that the above program works with a  $k \times 6$  matrix where  $k$  could be passed for `nrows`.

1. Why? because C does not care about the number of rows, only the number of cols.
2. And why is that? We'll have to understand 2-dim array addressing.

Example...

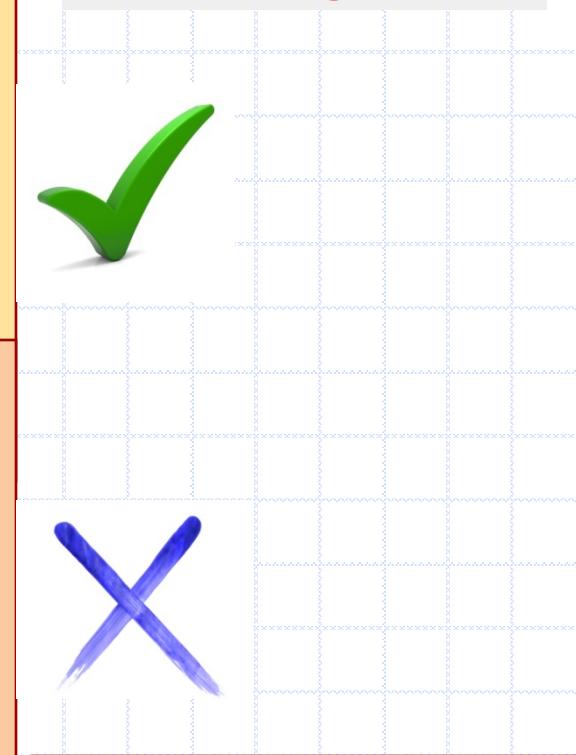
```
void marginals(double mat[ ][6], int nrows);
void main() {
    double mat[9][6];
    /* read the first 8 rows into mat */
    marginals(mat,8);
}
```

## Example calls for marginals

```
void marginals(double mat[ ][6], int nrows);
void main() {
    double mat[9][6];
    /* read 9 rows into mat */
    marginals(mat,10);
}
```

UNSAFE

The 10<sup>th</sup> row of mat[9][6] is not defined.  
So we may get a segmentation fault  
when marginals() processes the 10<sup>th</sup> row,  
i.e., i becomes 9.



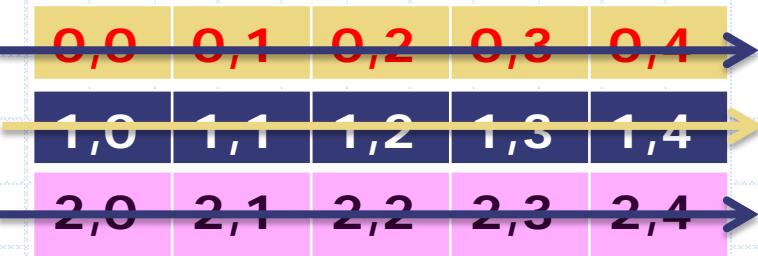
As with 1 dim arrays, allocate your array and stay within the limits allocated.

# Why is # of columns required?

- ◆ The **memory** of a computer is a **1D array!**
- ◆ 2D (or >2D) arrays are “**flattened**” into 1D to be stored in memory
- ◆ In C (and most other languages), arrays are flattened using **Row-Major** order
  - In case of 2D arrays, knowledge of number of columns is required to figure out where the next row starts.
  - **Last n-1 dimensions required for nD arrays**

# Row Major Layout

**mat[3][5]**



**Layout of mat[3][5] in memory**

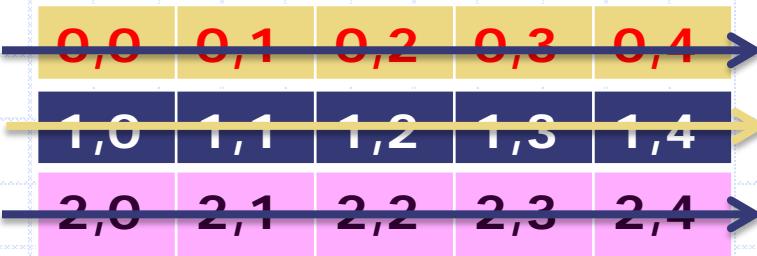


- for a 2D array declared as **mat[M][N]**, cell  $[i][j]$  is stored in memory at location  $i*N + j$  from start of mat.
- for k-D array  $arr[N_1][N_2]...[N_k]$ , cell  $[i_1][i_2]...[i_k]$  will be stored at location

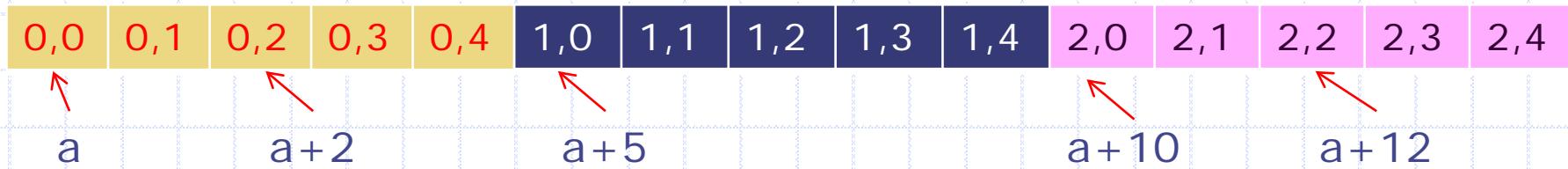
$$i_k + N_k * (i_{k-1} + N_{k-1} * (i_{k-2} + ( \dots + N_2 * i_1 ) \dots ))$$

# Row Major Layout

$\text{mat}[3][5]$



Layout of  $\text{mat}[3][5]$  in memory



- **About C implementation:**  $a = *mat$
- $*mat = \text{mat}[0]$ ,  $*(\text{mat}+1) = \text{mat}[1]$ ,  
 $*(\text{mat}+2) = \text{mat}[2], \dots$ . Each of which stores the reference to the corresponding row.
- That is, **mat** POINTS to the beginning of the array that stores the references to each of the rows.

# Array of Strings

◆ 2D array of char.

◆ Recall

- Strings are character arrays that end with a '\0'
- To display a string we can use printf with the %s placeholder.
- To input a string we can use scanf with %s.  
Only reads non-whitespace characters.

# Array of Strings: Example

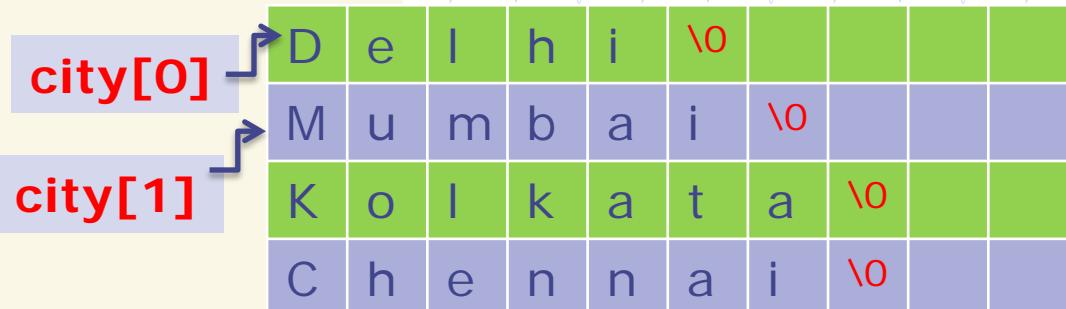
- ❖ Write a program that reads and displays the name of few cities of India

```
const int ncity = 4;
const int lencity = 10;

int main(){
    char city[ncity][lencity];
    int i;
    for (i=0; i<ncity; i++){
        scanf("%s", city[i]);
    }
    for (i=0; i<ncity; i++){
        printf("%d %s\n", i, city[i]);
    }
    return 0;
}
```

## INPUT

Delhi  
Mumbai  
Kolkata  
Chennai



## OUTPUT

- 0 Delhi
- 1 Mumbai
- 2 Kolkata
- 3 Chennai

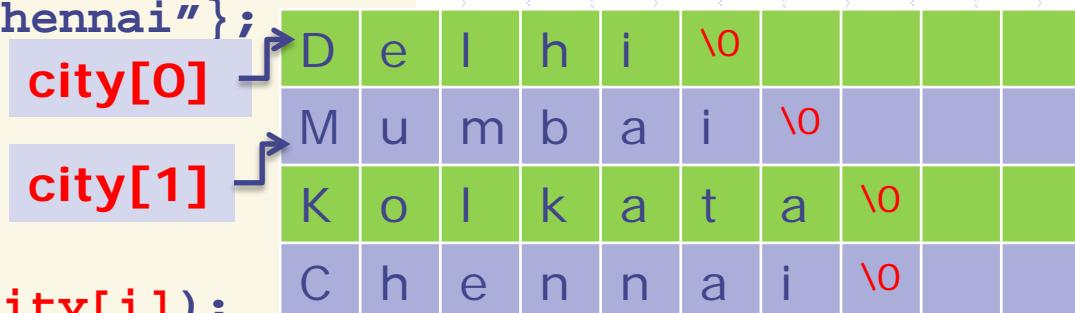
# Array of Strings: Example

◆ List initialization is also allowed:

```
const int ncity = 4;
const int lencity = 10;
```

```
int main(){
    char city[][lencity] = {"Delhi",
                            "Mumbai", "Kolkata", "Chennai"};
    int i;

    for (i=0; i<ncity; i++){
        printf("%d %s\n", i, city[i]);
    }
    return 0;
}
```

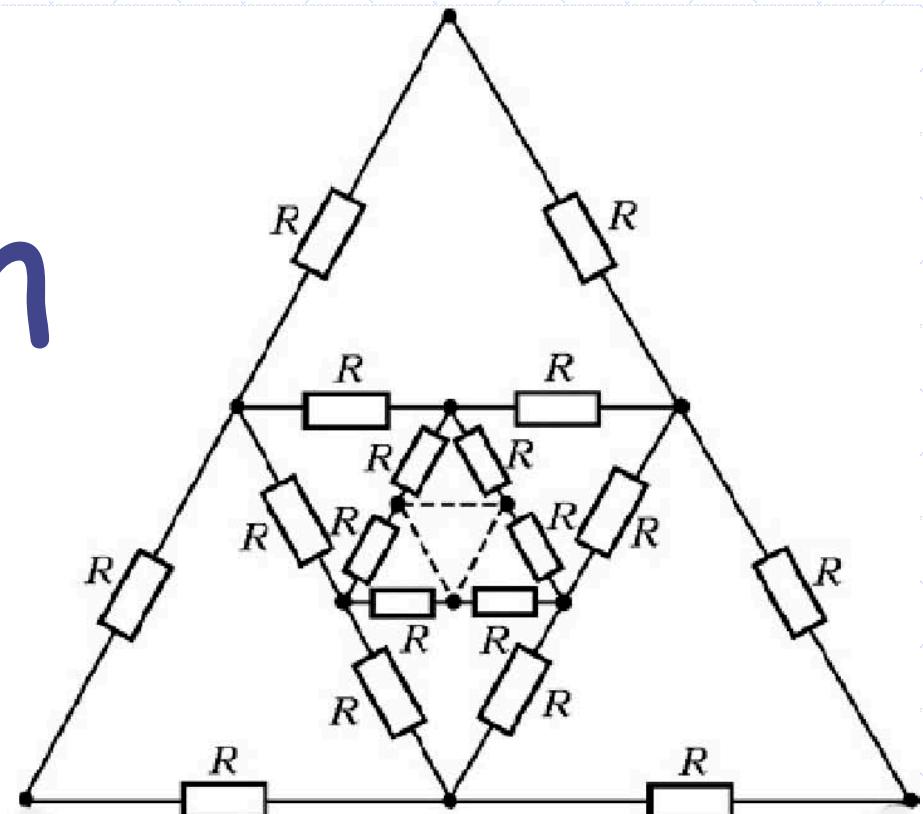
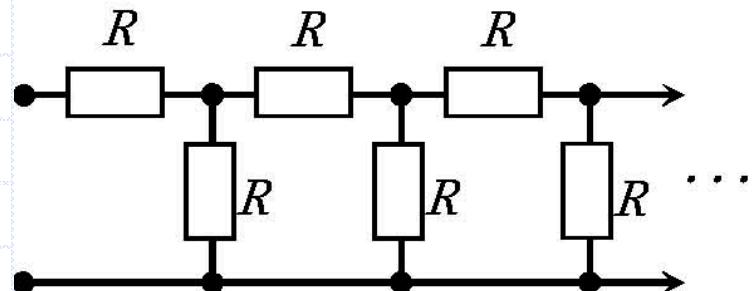


## OUTPUT

```
0 Delhi
1 Mumbai
2 Kolkata
3 Chennai
```

# ESC101: Introduction to Computing

## Recursion



# Recursion

◆ A function calling itself, directly or indirectly, is called a recursive function.

- The phenomenon itself is called recursion

◆ Examples:

- Factorial:

$$0! = 1$$
$$n! = n * (n-1)!$$

- Even and Odd:

$$\text{Even}(n) = (n == 0) \text{ || } \text{Odd}(n-1)$$
$$\text{Odd}(n) = (n != 0) \text{ && } \text{Even}(n-1)$$