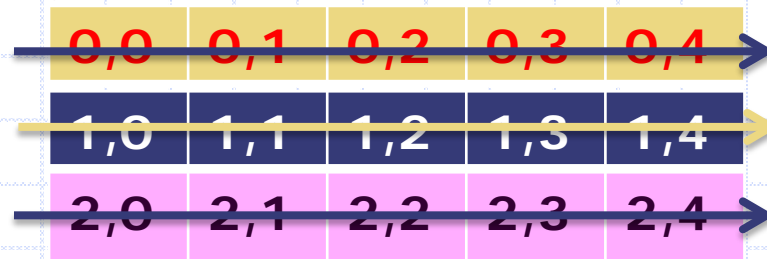


# Row Major Layout

`mat[3][5]`



Layout of `mat[3][5]` in memory

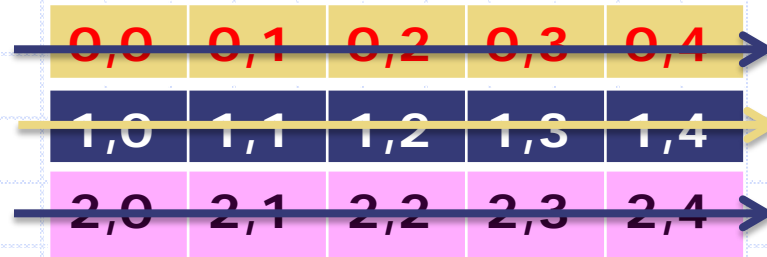


- for a 2D array declared as `mat[M][N]`, cell `[i][j]` is stored in memory at location  $i*N + j$  from start of `mat`.
- for  $k$ -D array `arr[N1][N2]...[Nk]`, cell `[i1][i2]...[ik]` will be stored at location

$$i_k + N_k * (i_{k-1} + N_{k-1} * (i_{k-2} + ( \dots + N_2 * i_1 ) \dots ))$$

# Row Major Layout

`mat[3][5]`



Layout of `mat[3][5]` in memory



- **About C interpretation:** `a = *mat`
- `*mat = mat[0]`, `*(mat+1) = mat[1]`,  
`*(mat+2) = mat[2]`,..... Each of which “stores” the reference to the corresponding row.
- That is, `mat` “points” to the beginning of the array that stores the references to each of the rows.

# Array of Strings: Example

- ◆ Write a program that reads and displays the name of few cities of India

```
const int ncity = 4;
const int lencity = 10;

int main(){
    char city[ncity][lencity];
    int i;

    for (i=0; i<ncity; i++){
        scanf("%s", city[i]);
    }

    for (i=0; i<ncity; i++){
        printf("%d %s\n", i, city[i]);
    }
    return 0;
}
```

## INPUT

Delhi  
Mumbai  
Kolkata  
Chennai

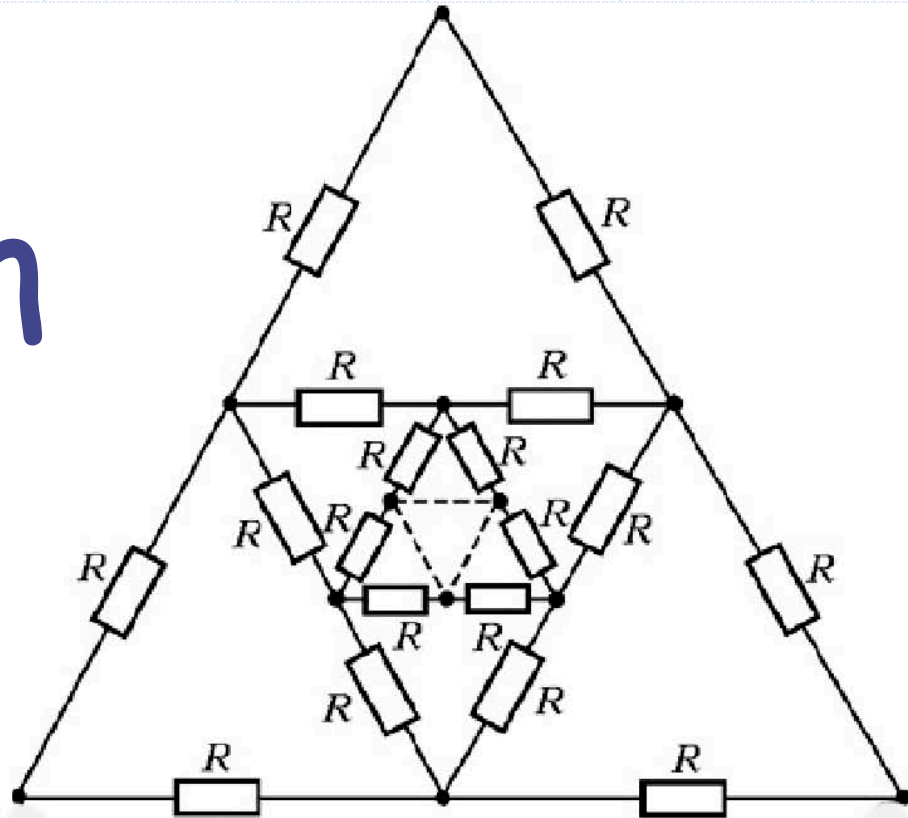
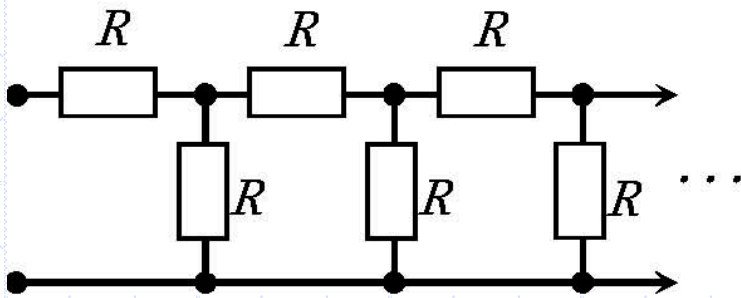
<b>city[0]</b>	D	e	l	h	i	\0			
<b>city[1]</b>	M	u	m	b	a	i	\0		
	K	o	l	k	a	t	a	\0	
	C	h	e	n	n	a	i	\0	

## OUTPUT

0 Delhi  
1 Mumbai  
2 Kolkata  
3 Chennai

# ESC101: Introduction to Computing

## Recursion



# Recursion

◆ A function calling itself, *directly* or *indirectly*, is called a *recursive function*.

- The phenomenon itself is called recursion

◆ Examples:

- Factorial:

$$0! = 1$$
$$n! = n * (n-1)!$$

- Even and Odd:

$$\text{Even}(n) = (n == 0) \ || \ \text{Odd}(n-1)$$
$$\text{Odd}(n) = (n != 0) \ \&\& \ \text{Even}(n-1)$$

# Recursive Functions: Properties

- ◆ The arguments **change** between the recursive calls

$$5! = 5 * 4! = 5 * 4 * 3! = \dots$$

- ◆ Change is towards a case for which solution is **known (base case)**
- ◆ There must be one or more **base cases**

$$0! \text{ is } 1$$

**Odd(0) is false**

**Even(0) is true**

# Recursion and Induction

*When programming recursively,  
think inductively*

- ◆ Mathematical induction for the natural numbers
- ◆ Structural induction for other recursively-defined types (to be covered later!)

# Recursion and Induction

When writing a recursive function,

- ◆ Write down a clear, concise *specification* of its behavior,
- ◆ Give an *inductive proof* that your code satisfies the specification.



# Constructing Recursive functions: Examples

Write a function `search(int a[], int n, int key)` that performs a sequential search of the array `a[0..n-1]` of `int`. Returns 1 if the key is found, otherwise returns 0.

How should we start? We have to think of the function `search()` in terms of search applied to a smaller array. Don't think in terms of loops...think recursion.

Here's a possibility ....

## search(a, n, key)

**Base case:** If  $n$  is 0, then, return 0.

**Otherwise:** /\*  $n > 0$  \*/

1. compare last item,  $a[n-1]$ , with key.
2. if  $a[n-1] == \text{key}$ , return 1.
3. search in array  $a$ , up to size  $n-1$ .
4. return the result of this "smaller" search.

**a**      search(a, 10, 3)

31	4	10	35	59	31	3	25	35	11
----	---	----	----	----	----	---	----	----	----

Either 3 is  $a[9]$ ; or search(a, 10, 3) is same as the result of search for 3 in the array starting at  $a$  and of size 9.

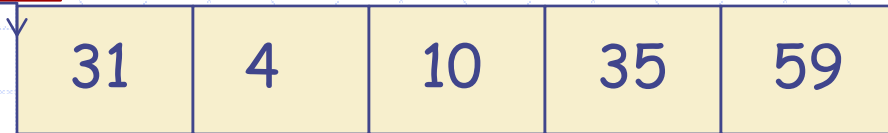
```

1. int search(int a[], int n, int key) {
2.   if (n==0) return 0;
3.   if (a[n-1] == key) return 1;
4.   return search(a,n-1,key);
5. }

```

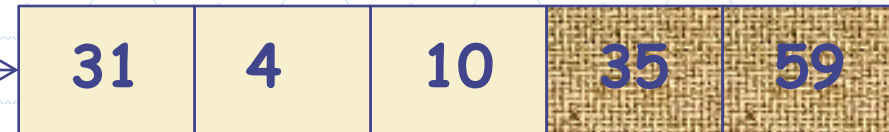
Let us do a quick trace.

**a** E.g., (0) search(a,5,10)



a[4] is 59, not 10. call search(a,4,10)

**a** (2) search(a,3,10)



a[2] is 10, return 1

**a** (1) search(a,4,10)



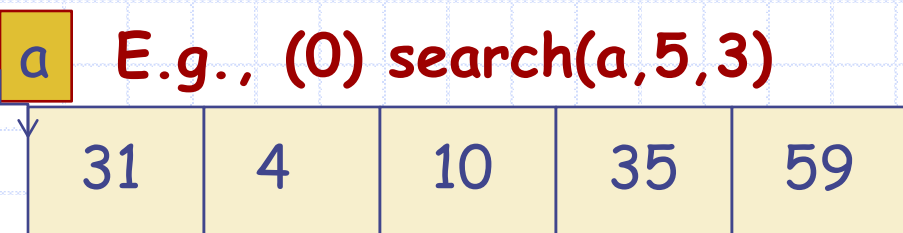
a[3] is 35, calls search(a,3,10)

```

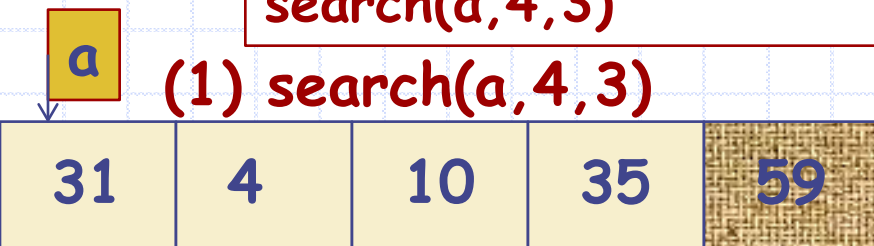
1. int search(int a[], int n, int key) {
2.   if (n==0) return 0;
3.   if (a[n-1] == key) return 1;
4.   return search(a,n-1,key);
5. }

```

Let us do another quick trace.



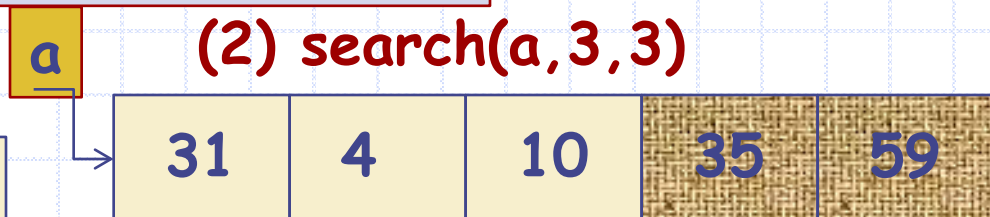
a[4] is 59, not 3. call search(a,4,3)



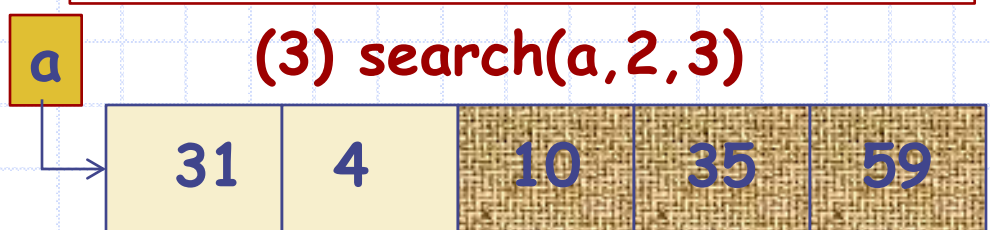
a[3] is 35, calls search(a,3,3)



**a** (5) search(a,0,3) returns 0



a[2] is 10, calls search(a,2,3)



a[1] is 4, calls search(a,1,3)



a[0] is 31, calls search(a,0,3)

```

1. int search(int a[], int n, int key) {
2.   if (n==0) return 0;
3.   if (a[n-1] == key) return 1;
4.   return search(a,n-1,key);
5. }

```

a	search(a,5,3)	
↓	31	4
	10	
	35	59

	function	called by	return address	return value
	search(a,5,3)	main()	---	
	search(a,4,3)	search(a,5,3)	search.5	
	search(a,3,3)	search(a,4,3)	search.4	
	search(a,2,3)	search(a,3,3)	search.3	
	search(a,1,3)	search(a,2,3)	search.2	
	search(a,0,3)	search(a,1,3)	search.1	

Stack



recursion exits here

A state of the stack

```

1. int search(int a[], int n, int key) {
2.   if (n==0) return 0;
3.   if (a[n-1] == key) return 1;
4.   return search(a,n-1,key);
5. }

```

a	search(a,5,3)	
↓	31	4
	10	
	35	59

	function	called by	return address	return value
	search(a,5,3)	main()	---	
	search(a,4,3)	search(a,5,3)	search.5	
	search(a,3,3)	search(a,4,3)	search.4	
	search(a,2,3)	search(a,3,3)	search.3	
	search(a,1,3)	search(a,2,3)	search.2	
	search(a,0,3)	search(a,1,3)	search.1	<b>0</b>

Stack



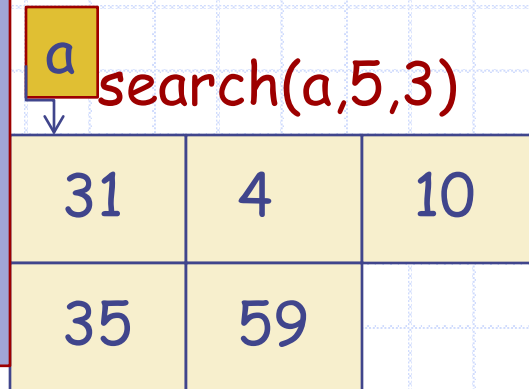
recursion exits here

A state of the stack

```

1. int search(int a[], int n, int key) {
2.   if (n==0) return 0;
3.   if (a[n-1] == key) return 1;
4.   return search(a,n-1,key);
5. }

```



Stack  
↓

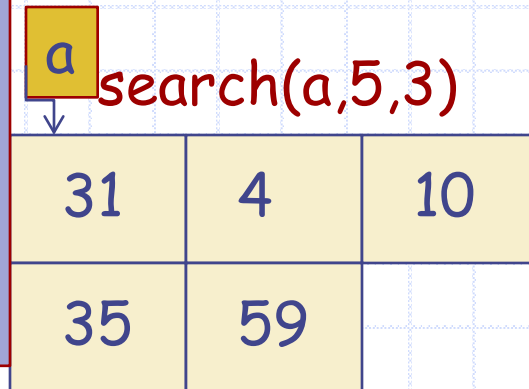
function	called by	return address	return value
search(a,5,3)	main()	---	
search(a,4,3)	search(a,5,3)	search.5	
search(a,3,3)	search(a,4,3)	search.4	
search(a,2,3)	search(a,3,3)	search.3	
search(a,1,3)	search(a,2,3)	search.2	0

state of the stack

```

1. int search(int a[], int n, int key) {
2.   if (n==0) return 0;
3.   if (a[n-1] == key) return 1;
4.   return search(a,n-1,key);
5. }

```



	function	called by	return address	return value
	search(a,5,3)	main()	---	
	search(a,4,3)	search(a,5,3)	search.5	
	search(a,3,3)	search(a,4,3)	search.4	
Stack ↓	search(a,2,3)	search(a,3,3)	search.3	0

A state of the stack



```

1. int search(int a[], int n, int key) {
2.   if (n==0) return 0;
3.   if (a[n-1] == key) return 1;
4.   return search(a,n-1,key);
5. }

```

**a** search(a,5,3)

31	4	10
35	59	

	function	called by	return address	return value
Stack ↓	search(a,5,3)	main()	---	
	search(a,4,3)	search(a,5,3)	search.5	0

A state of the stack

```

1. int search(int a[], int n, int key) {
2.   if (n==0) return 0;
3.   if (a[n-1] == key) return 1;
4.   return search(a,n-1,key);
5. }

```

**a** search(a,5,3)

31	4	10
35	59	

Stack

function	called by	return address	return value
search(a,5,3)	main()	---	0

A state of the stack

```
1. int search(int a[], int n, int key) {  
2.   if (n==0) return 0;  
3.   if (a[n-1] == key) return 1;  
4.   return search(a,n-1,key);  
5. }
```

a search(a,5,3)  
↓

31	4	10
35	59	

search(a,5,3) returns 0. Recursion call stack terminates.

# Searching in an Array

- ◆ We can have other recursive formulations
- ◆ **Search1:** `search(a, start, end, key)`
  - Search key between `a[start]...a[end]`

if `start > end`, return 0;

if `a[start] == key`, return 1;

return `search(a, start+1, end, key)`;

# Searching in an Array

- ◆ One more recursive formulation
- ◆ **Search2**: search (a, start, end, key)
  - Search key between a[start]...a[end]

if start > end, return 0;

mid = (start + end)/2 ;

if a[mid]==key, return 1;

return search(a, start, mid-1, key)

|| search(a, mid+1, end, key);