

# ESC101: Introduction to Computing

# Sorting



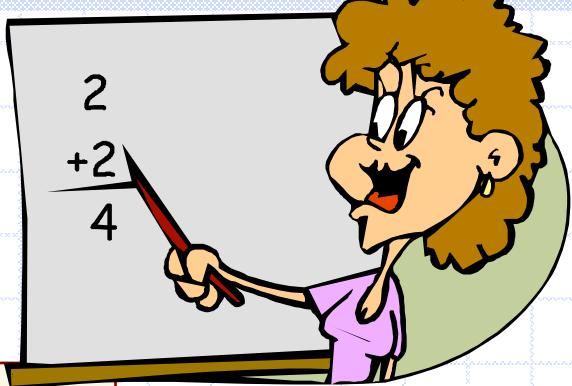
# Sorting

- ◆ Given a list of integers (in an array), arrange them in ascending order.
  - Or descending order

INPUT ARRAY	5	6	2	3	1	4
OUTPUT ARRAY	1	2	3	4	5	6

- ◆ Sorting is an extremely important problem in computer science.
  - A common problem in everyday life.
  - Example:
    - ◆ Contact list on your phone.
    - ◆ Ordering marks before assignment of grades.

# What's easy to do in a Sorted Array?



Clearly, searching for a key is fast.

Rank Queries: find the  $k^{\text{th}}$  largest/smallest value.  
Quantile: 90%ile—the key value in the array such that 10% of the numbers are larger than it.

40	50	55	60	70	75	80	85	90	92
----	----	----	----	----	----	----	----	----	----

Marks in an exam: sorted

90 percentile : 90  
80 percentile : 85  
10 percentile: 40  
50 percentile: 70  
(also called median)

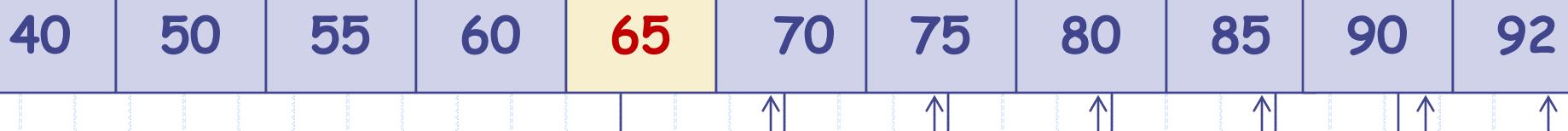
# Sorted array have difficulty with

- ◆ inserting a new element while preserving the sorted structure.
- ◆ deleting an existing element (while preserving the sorted structure).
- ◆ In both cases, there may be need to shift elements to the right or left of the index corresponding to insertion or deletion.



Example: Insert 65.

1. Find index where 65 needs to be inserted



2. Shift right from index 5 to create space.

3. Insert 65

May have to shift  $n-1$  elements in the worst case.

# Sorting

◆ Many well known sorting Algorithms

- Selection sort
- Merge sort
- Quick sort
- Bubble sort
- ...

◆ Special cases also exist for specific problems/data sets

◆ Different runtime

◆ Different memory requirements

# Selection Sort

- ◆ Select the largest element in your array and swap it with the first element of the array.
- ◆ Consider the sub-array from the second element to the last, as your current array and repeat Step 1.
- ◆ Stop when the array has only one element.
  - Base case, trivially sorted

# Selection Sort: Pseudo code

```
selection_sort(a, start, end) {
    if (start == end) /* base case, one elt => sorted */
        return;

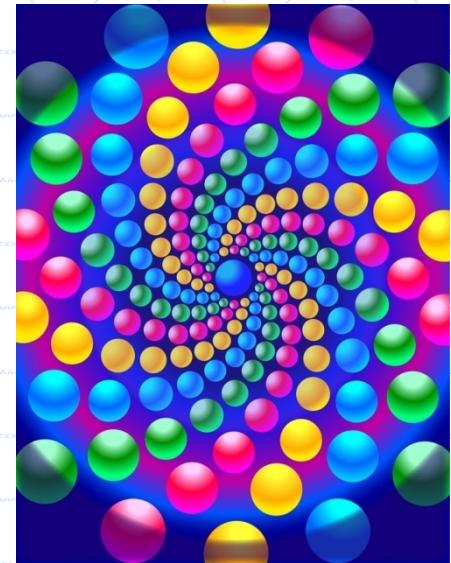
    idx_max = find_idx_of_max_elt(a, start, end);
    swap(a, idx_max, start);
    selection_sort(a, start+1, end);
}
```

```
swap(a, i, j) {
    tmp = a[i];
    a[i] = a[j];
    a[j] = tmp;
}
```

```
main() {
    arr[] = { 5, 6, 2, 3, 1, 4 };
    selection_sort(arr, 0, 5);
    /* print arr */
}
```

# Selection Sort: Properties

- ◆ Is the pseudo code iterative or recursive?
- ◆ What is the estimated run time when input array has  $n$  elements
  - for swap      Constant
  - for find\_idx\_of\_max\_elt       $\propto n$
  - for selection\_sort      On next slide
- ◆ Practice: Write C code for iterative and recursive versions of selection sort.



# Selection Sort: Time Estimate

◆ Recurrence

$$T(n) = T(n - 1) + k_1 \times n + k_2$$

◆ Solution

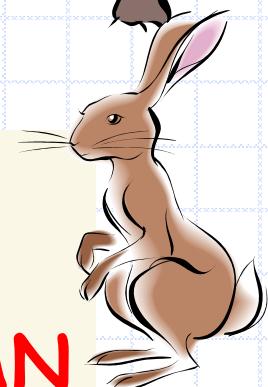
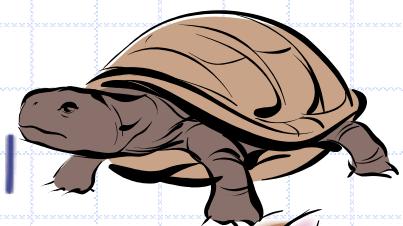
$$T(n) \propto n(n + 1)$$

◆ Or simply

$$T(n) \propto n^2$$

Selection sort runs in time proportional  
to the square of the size of the array  
to be sorted.

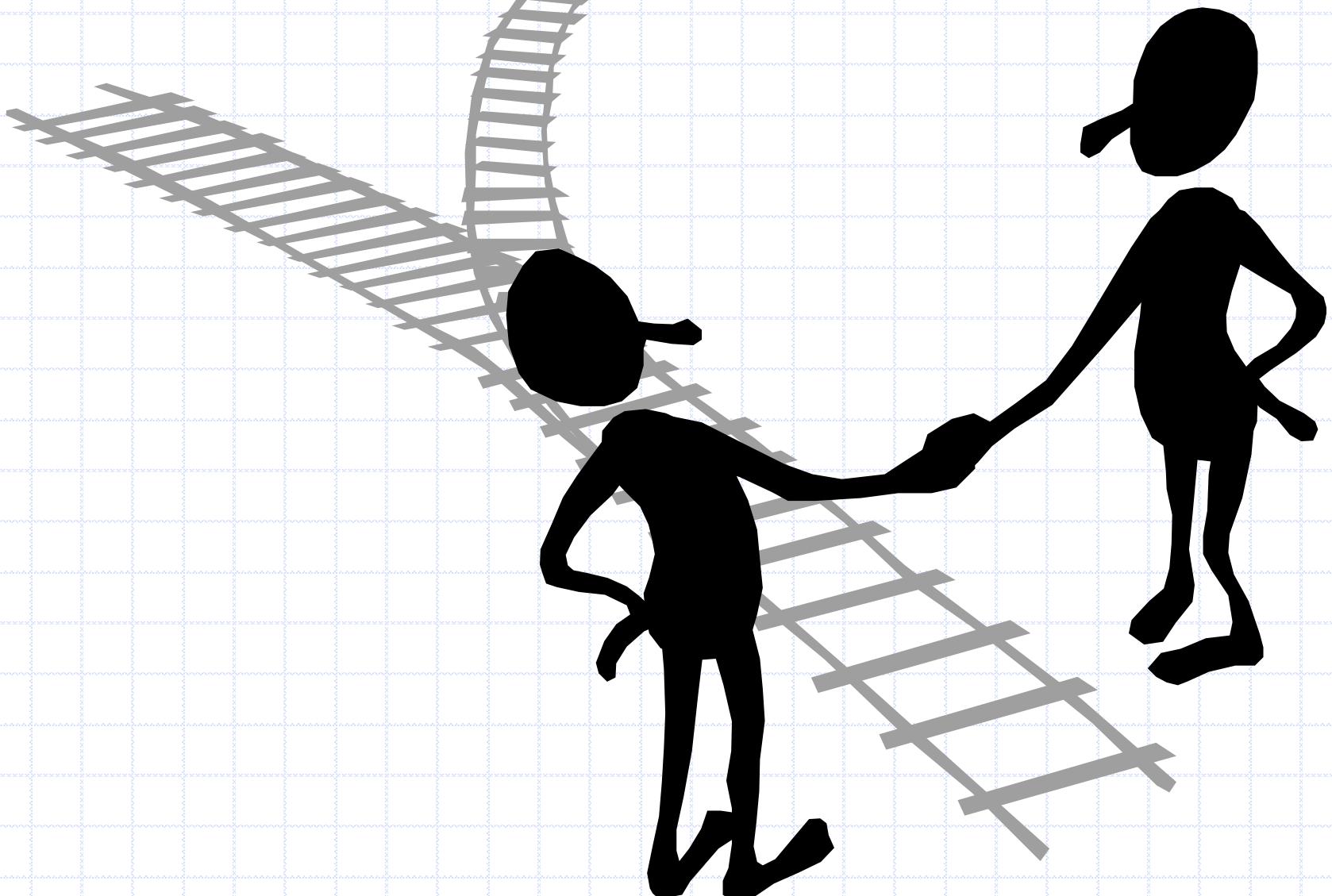
Can we do  
better?  
**YES WE CAN**



# Merging Two Sorted Arrays

- ◆ Input: Array A of size n & array B of size m.
- ◆ Create an empty array C of size  $n + m$ .
- ◆ Variables i , j and k
  - array variables for the arrays A, B and C resp.
- ◆ At each iteration
  - compare the  $i^{\text{th}}$  element of A (say u) with the  $j^{\text{th}}$  element of B (say v)
  - if u is smaller, copy u to C; increment i and k,
  - otherwise, copy v to C; increment j and k,

# Merging Two Sorted Arrays



# Time Estimate

- ◆ Number of steps  $\propto 3(n + m)$ .
  - The constant 3 is not very important as it does not vary with different sized arrays.
- ◆ Now suppose A and B are halves of an array of size n (both have size  $n/2$ ).
- ◆ Number of steps =  $3n$ .

$$T(n) \propto n$$

# MergeSort

- ◆ Merge function can be used to sort an array
  - recursively!
- ◆ Given an array  $C$  of size  $n$  to sort
  - Divide it into Arrays A and B of size  $n/2$  each (approx.)
  - Sort A into  $A'$  using MergeSort
  - Sort B into  $B'$  using MergeSort
  - Merge  $A'$  and  $B'$  to give  $C' \equiv C$  sorted
- ◆ Can we reduce #of extra arrays ( $A'$ ,  $B'$ ,  $C'$ )?

Recursive calls.  
Base case?

$n \leq 1$

```
/*Sort ar[start, ..., start+n-1] in place */
void merge_sort(int ar[], int start, int n) {
    if (n>1) {
        int half = n/2;
        merge_sort(ar, start, half);
        merge_sort(ar, start+half, n-half);
        merge(ar, start, n);
    }
}
```

```
int main() {
    int arr[]={2,5,4,8,6,9,8,6,1,4,7};
    merge_sort(arr,0,11);
    /* print array */
    return 0;
}
```

```
void merge(int ar[], int start, int n) {  
    int temp[MAX_SZ], k, i=start, j=start+n/2;  
    int lim_i = start+n/2, lim_j = start+n;  
    for(k=0; k<n; k++) {  
        if ((i < lim_i) && (j < lim_j)) {// both active  
            if (ar[i] <= ar[j]) { temp[k] = ar[i]; i++; }  
            else { temp[k] = ar[j]; j++; }  
        } else if (i == lim_i) // 1st half done  
        { temp[k] = ar[j]; j++; } // copy 2nd half  
        else // 2nd half done  
        { temp[k] = ar[i]; i++; } // copy 1st half  
    }  
    for (k=0; k<n; k++)  
        ar[start+k]=temp[k]; // in-place  
}
```

# Time Estimate

```
void merge_sort(int a[], int s, int n) { T(n)
    if (n>1) {
        int h = n/2;
        merge_sort(a, s, h); T(n/2)
        merge_sort(a, s+h, n-h); T(n-n/2)≈T(n/2)
        merge(a, s, n); ≈ 4n
    }
}
```

# Time Estimate

$$T(n) = 2T(n/2) + 4n$$

$$= 2(2T(n/4) + 4n/2) + 4n = 2^2T(n/4) + 8n$$

$$= 2^2(2T(n/8) + 4n/4) + 4n = 2^3T(n/8) + 12n$$

= ... // keep going for k steps

$$= 2^k T(n/2^k) + k*4n$$

Assume  $n = 2^k$  for some k. Then,

$$T(n) = n*T(1) + 4n*\log_2 n$$

$$T(n) \propto n \log_2 n$$

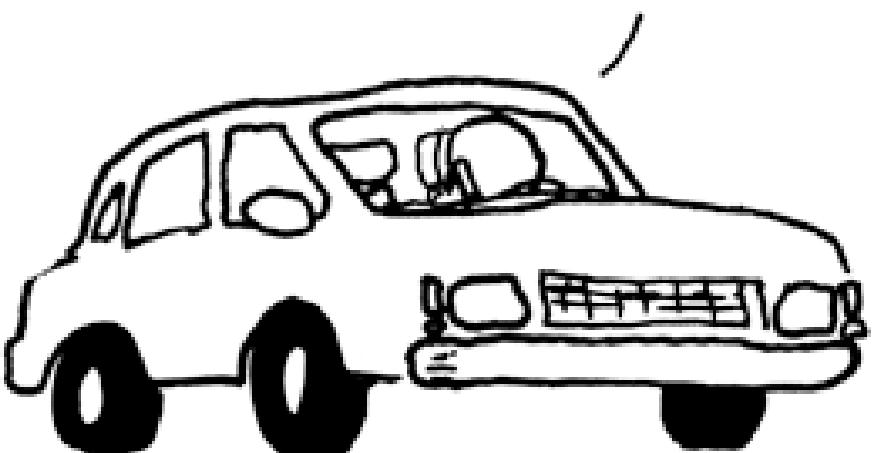
# Order Notation

Why worry  
about  $O(n)$  vs  
 $O(n^2)$  vs  $O(\dots)$   
algo?

I'M JUST OUTSIDE TOWN, SO I SHOULD  
BE THERE IN FIFTEEN MINUTES.

ACTUALLY, IT'S LOOKING  
MORE LIKE SIX DAYS.  
I

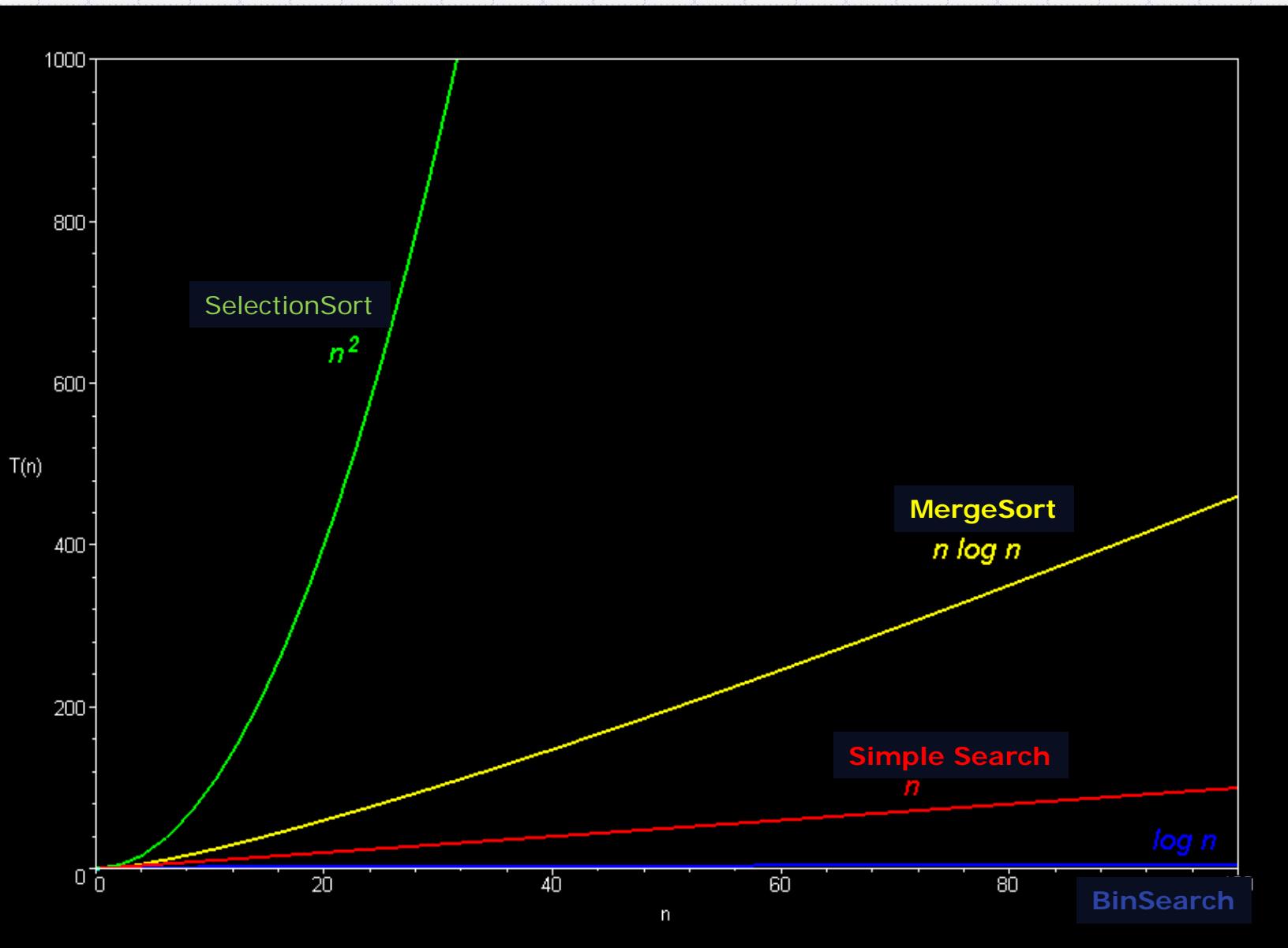
NO, WAIT, THIRTY SECONDS.



THE AUTHOR OF THE WINDOWS FILE  
COPY DIALOG VISITS SOME FRIENDS.

<http://xkcd.com/612/>

# Time Estimates...



Source:  
<http://science.slc.edu/~jmarshall/courses/2002/spring/cs50/BigO/>

# QuickSort-- Partition Routine

A useful sub-routine (function) for many problems, including quicksort(), the most popularly used sorting method.

1. Partition takes an array  $a[]$  of size  $n$  and a value called the pivot.
2. The pivot is an element in the array is usually chosen as  $a[0]$ .
3. Partition re-arranges the array elements into two parts:
  - a) the left part has all elements  $\leq$  pivot.
  - b) the right part has all elements  $\geq$  pivot.
4. Partition returns the index of the beginning of the right part.

Let us see an example.

1. **Partition** takes an array  $a[]$  of size  $n$  and a value called the pivot.
2. The pivot is an element in the array is usually chosen as  $a[0]$ .
3. Partition re-arranges the array elements into two parts:
  - a) all elements in the left part are  $\leq$  pivot
  - b) all elements in the right part are  $\geq$  pivot

Input Array  $a[]$ , size is  $n : 11$

31	4	10	35	59	31	3	25	35	11	0
----	---	----	----	----	----	---	----	----	----	---

Call to  $\text{partition}(a, 11)$ . Pivot element is assumed to be  $a[0]: 31$

0	4	10	11	25	3	31	31	35	35	31
---	---	----	----	----	---	----	----	----	----	----

left partition

right partition

return value is 6

# COMMENTS

Multiple “partitions” of an array are possible, even for the same pivot. They all would satisfy the above specification.

Note: Partition **DOES NOT** sort the array. It is “weaker” than sorting. But it is useful step towards sorting (useful for other problems also).

1. **partition(int a[], int n)**. pivot will be  $a[0]$ .
2. Partition re-arranges the array elements into two parts:
  - a) the left part has all elements  $\leq$  pivot
  - b) the right part has all elements  $\geq$  pivot
3. Partition should return either the first index of the right part or the last index of the left part. (Both answers would be acceptable).

Designing partition: Goal is to have linear time complexity, meaning that the number of comparisons and exchanges of items must be linear in the size.

Also, we will do partition *in place* - that is, without using extra arrays.

Can you do it easily if you have extra arrays?

a

## Designing Partition

31 4 10 35 59 31 3 25 35 11 0

0

31

10

l

pivot

r

1. Keep two integer variables denoting indices: l starts at the left end and r starts at the right end.
2. pivot is  $a[0]$  which is 31.
3. Value of pivot will not change during partition.

### Basic Step in Partition Loop:

As long as  $a[l] \leq \text{pivot}$ , advance l by 1.

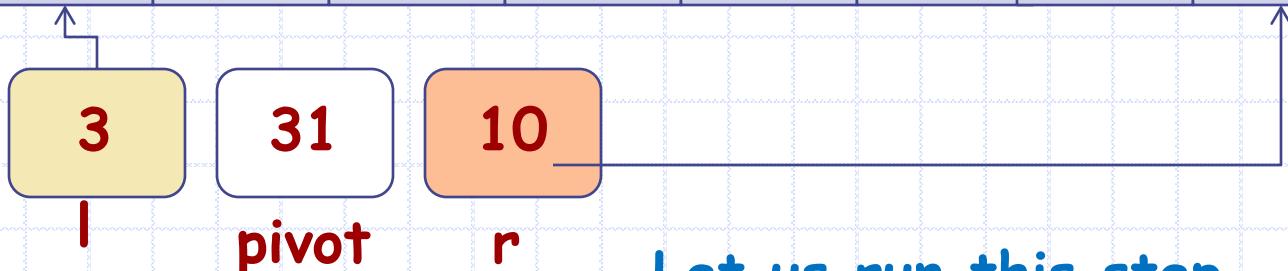
As long as  $a[r] \geq \text{pivot}$ , decrease r by 1.

If  $l < r$ , Exchange  $a[l]$  with  $a[r]$ .  
advance l by 1; decrement r by 1

a

## Designing Partition

31 4 10 35 59 31 3 25 35 11 0



Basic Step in Partition Loop:

As long as  $a[l] \leq \text{pivot}$ ,  
advance l by 1.

As long as  $a[r] \geq \text{pivot}$ ,  
decrease r by 1.

If  $l < r$ , Swap  $a[l]$  with  
 $a[r]$ .  
advance l by 1;  
decrement r by 1

Let us run this step  
on the above array

1. First loop terminates, with l as 3.
2. Second loop terminates immediately, with r as 10.

Now we swap  $a[3]$  with  $a[10]$

a

## Designing Partition

31 4 10 0 59 31 3 25 35 11 35

4

31

9

pivot

r

Basic Step in Partition Loop:

As long as  $a[l] < \text{pivot}$ ,  
advance l by 1.

As long as  $a[r] > \text{pivot}$ ,  
decrease r by 1.

Swap  $a[l]$  with  $a[r]$ .  
advance l by 1;  
decrement r by 1

Let us run this step  
on the above array

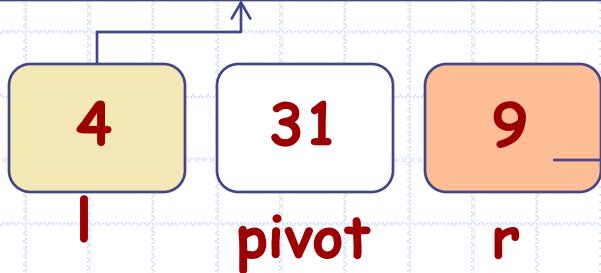
Swap and Advance

1. swap  $a[3]$  with  $a[10]$
2. Advance l to 4
3. Decrement r to 9

a

## Designing Partition

31 4 10 0 59 31 3 25 35 11 35



Basic Step in Partition Loop:

As long as  $a[l] < \text{pivot}$ ,  
advance l by 1.

As long as  $a[r] > \text{pivot}$ ,  
decrease r by 1.

Swap  $a[l]$  with  $a[r]$ .  
advance l by 1;  
decrement r by 1

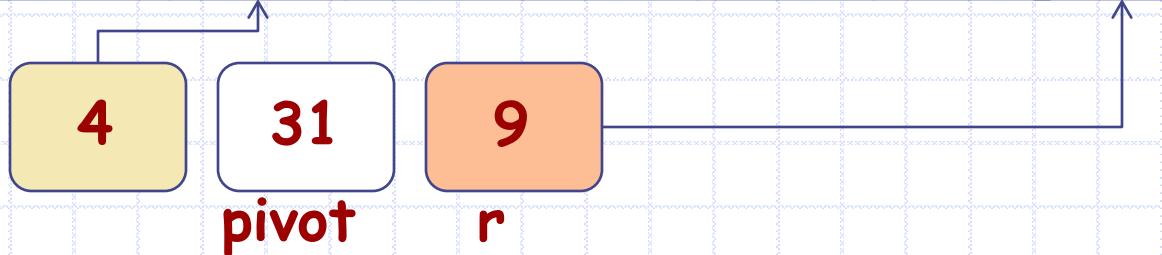
Let us run this step  
on the above array

Invariant

1.  $a[0] \dots a[l-1]$  are all  $\leq$  pivot.
2.  $a[r+1] \dots a[n-1]$  are all  $\geq$  pivot.

# Designing Partition

a	31	4	10	0	59	31	3	25	35	11	35
---	----	---	----	---	----	----	---	----	----	----	----



## Basic Step in Partition Loop:

As long as  $a[l] \leq \text{pivot}$ ,  
advance  $l$  by 1.

As long as  $a[r] \geq \text{pivot}$ ,  
decrease  $r$  by 1.

If  $l < r$ , Swap  $a[l]$  with  
 $a[r]$ .  
advance  $l$  by 1;  
decrement  $r$  by 1

Now what should we do?  
We can run the basic step again.

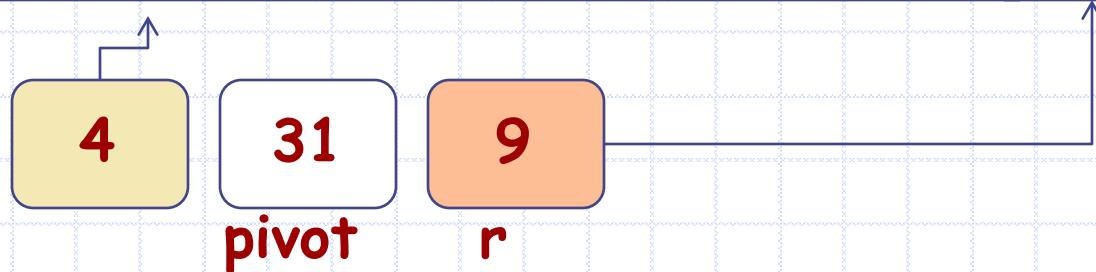
## Loop for $l$

1. The loop for  $l$  terminates immediately since  $59 > 31$ .

a

## Designing Partition

31 4 10 0 59 31 3 25 35 11 35



**Basic Step in Partition Loop:**

As long as  $a[l] \leq \text{pivot}$ ,  
advance  $l$  by 1.

As long as  $a[r] \geq \text{pivot}$ ,  
decrease  $r$  by 1.

If  $l < r$ , Swap  $a[l]$  with  
 $a[r]$ .  
advance  $l$  by 1;  
decrement  $r$  by 1

Now what should we do?  
We can run the basic step again.

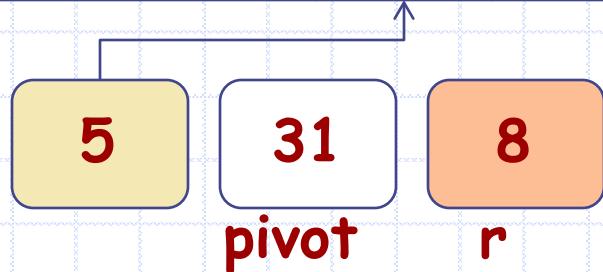
**Swap and advance**

1. Swap 59 with 11
2. Increment  $l$  by 1
3. Decrement  $r$  by 1

a

## Designing Partition

31	4	10	0	11	31	3	25	35	59	35
----	---	----	---	----	----	---	----	----	----	----



Basic Step in Partition Loop:

As long as  $a[l] \leq \text{pivot}$ ,  
advance  $l$  by 1.

As long as  $a[r] \geq \text{pivot}$ ,  
decrease  $r$  by 1.

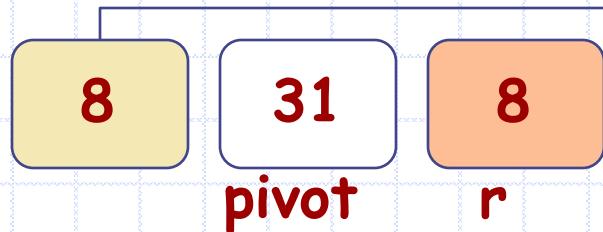
If  $l < r$ , Swap  $a[l]$  with  $a[r]$ .  
advance  $l$  by 1;  
decrement  $r$  by 1

Now what should we do?  
We can run the basic step again.

a

## Designing Partition

31 4 10 0 11 31 3 25 35 59 35



### Basic Step in Partition Loop:

As long as  $a[l] \leq \text{pivot}$ ,  
advance  $l$  by 1.

As long as  $a[r] \geq \text{pivot}$ ,  
decrease  $r$  by 1.

If  $l < r$ , Swap  $a[l]$  with  
 $a[r]$ .  
advance  $l$  by 1;  
decrement  $r$  by 1

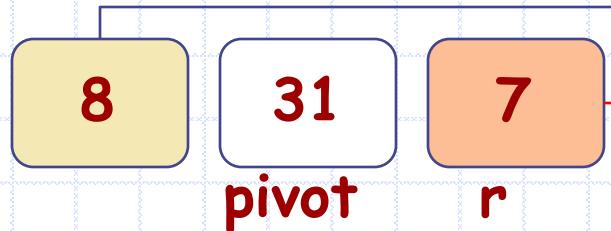
### Loop for $l$

1. The loop for  $l$  terminates at  $l=8$ .

a

## Designing Partition

31 4 10 0 11 31 3 25 35 59 35



### Basic Step in Partition Loop:

As long as  $a[l] \leq \text{pivot}$ ,  
advance l by 1.

As long as  $a[r] \geq \text{pivot}$ ,  
decrease r by 1.

If  $l < r$ , Swap  $a[l]$  with  
 $a[r]$ .  
advance l by 1;  
decrement r by 1

### Loop for r

1. The loop for l terminates  
at  $r=7$ .

## Designing Partition

31 4 10 0 11 31 3 25 35 59 35



Basic Step in Partition Loop:

Partition is almost done!

As long as  $a[l] < \text{pivot}$ ,  
advance  $l$  by 1.

As long as  $a[r] > \text{pivot}$ ,  
decrease  $r$  by 1.

Swap  $a[l]$  with  $a[r]$ .  
advance  $l$  by 1;  
decrement  $r$  by 1

Why? From invariant:

1. What should we return?
2. We can return  $r$ .

```
int partition(int a[], int n) {
    int l = 0, r = n-1, pivot = a[0];
    while (l <= n-1 && r >= 0) {
        while (a[l] <= pivot) { l=l+1; }
        while (a[r] >= pivot) { r=r-1; }
        if(l<r) {
            swap(a, l, r);
            l = l+1; r = r-1;
        } else {
            /* move pivot to its position */
            swap(a, l-1, 0);
            return l-1;
        }
    }
}
```

# The Partition function



We designed a function `int partition(int a[], int n)` that returns an index `pindex` of the array `a[]` such that for any `a[]` with `n >= 2`, all the following are true.

1. `pindex` lies between 0 and `n-2`,
  2. all items in `a[0..pindex]` are  $\leq$  pivot,
  3. all items in `a[pindex+1...n-1]` are  $\geq$  pivot,
  4. Number of operations required by partition is  $O(n)$ , that is bounded by  $c \cdot n$  for some constant  $c$ .
- Required only a single pass over the array: each element is touched once.

# Pivoting choices

Pivot may be chosen to be any value of  $a[]$ . Some choices are

1. Pivot is  $a[0]$ : simple choice.
2. Pivot is some random member of  $a[]$ : randomized pivot choice.
3. Pivot is the median element of  $a[]$ . This gives the most equal sized partitions, but is much more complicated.

Suppose we wish to sort the array  $a[]$ .

After the call  $pindex = \text{partition}(a, n)$

1. each element of  $a[0..pindex-1] \leq \text{pivot}$ .
2. each element of  $a[pindex..n-1] \geq \text{pivot}$ .

So after the call to  $\text{partition}()$ , to sort  $a[]$ , we can just

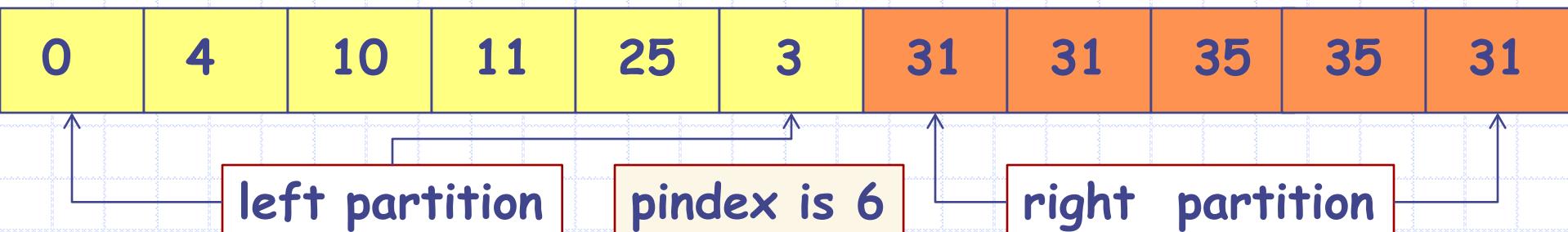
1. sort the array  $a[0..pindex-1]$ , and,
2. sort the array  $a[pindex..n-1]$ .

For example, consider the array.

Input Array  $a[]$ , size is  $n : 11$

31	4	10	35	59	31	3	25	35	11	0
----	---	----	----	----	----	---	----	----	----	---

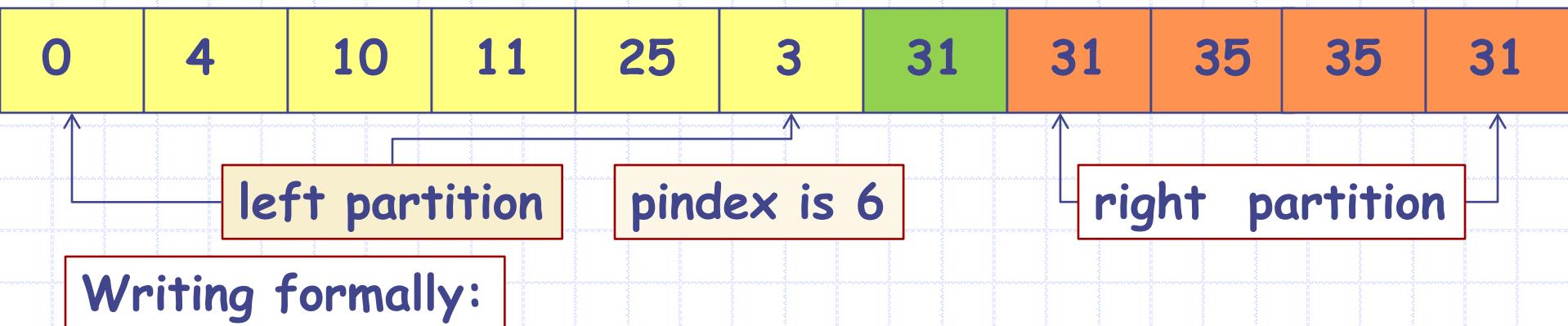
After call to `partition(a, 11)`. Pivot element is assumed to be `a[0]`: 31



Obvious from diagram: to sort `a[]`, we can sort the left partition and the right partition independently.

How should we do the sorting: Any way we wish, but... how about choosing the same algorithm, that is, run `partition` on each half again (and then again on smaller parts—this is recursion)

After call to `partition(a, 11)`. Pivot element is assumed to be `a[0]`: 31



```
void qsort(int a[], int n) {  
    int pindex;  
    if (n<=1) return; /* nothing to sort */  
    else {  
        pindex = partition(a,n);  
        qsort(a,pindex);  
        qsort(a+pindex+1, n-pindex-1);  
    }  
}
```

This is a recursive program

These are recursive calls.

31 4 10 35 59 31 3 25 35 11 31



Let us now run qsort on the above array: n is 11.

STACK

Function called

Called by function

n

pindex

Return address

main()

---

---

---

---

qsort(a, 11)

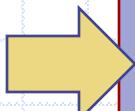
main()

11

-

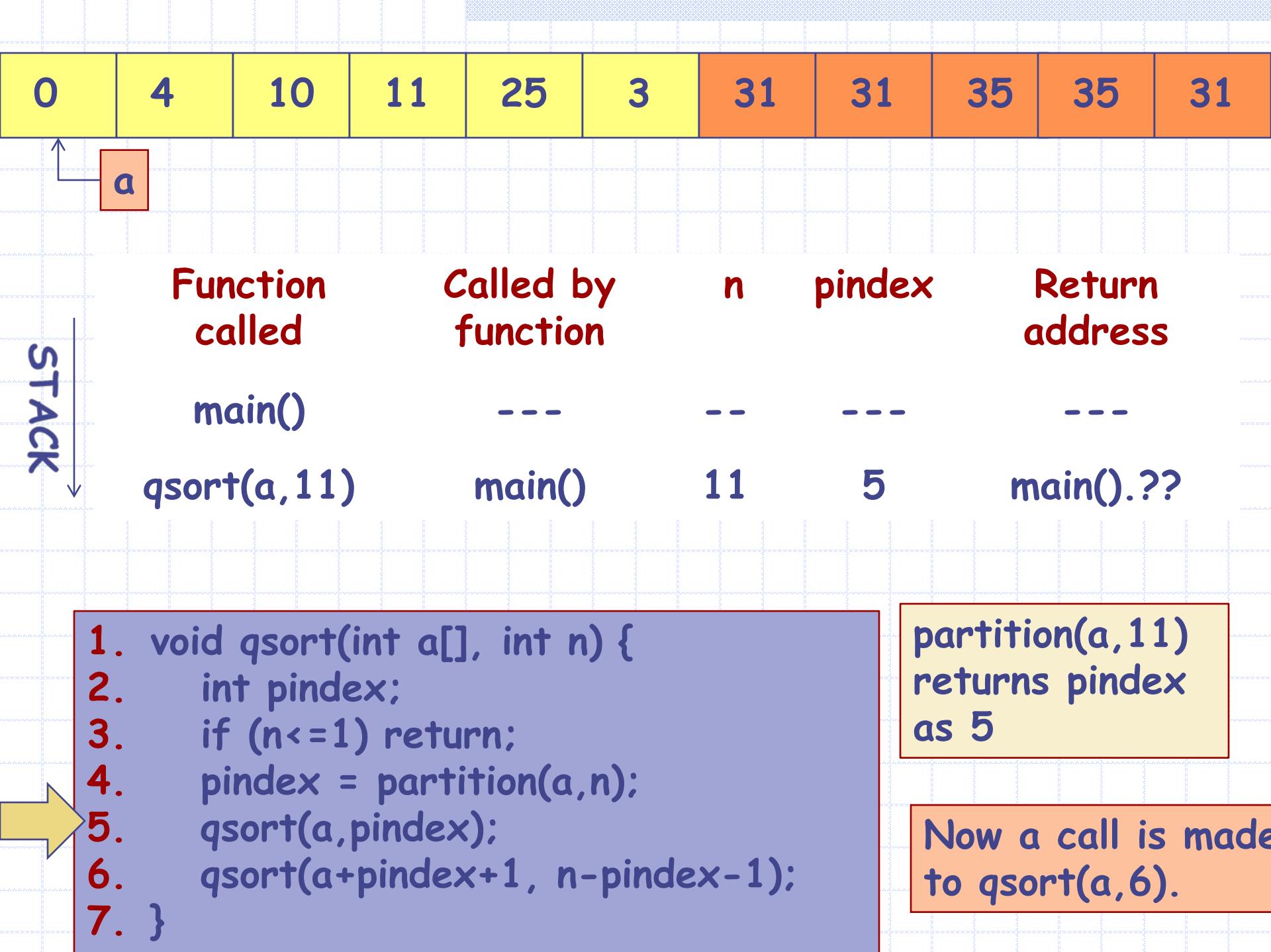
main().??

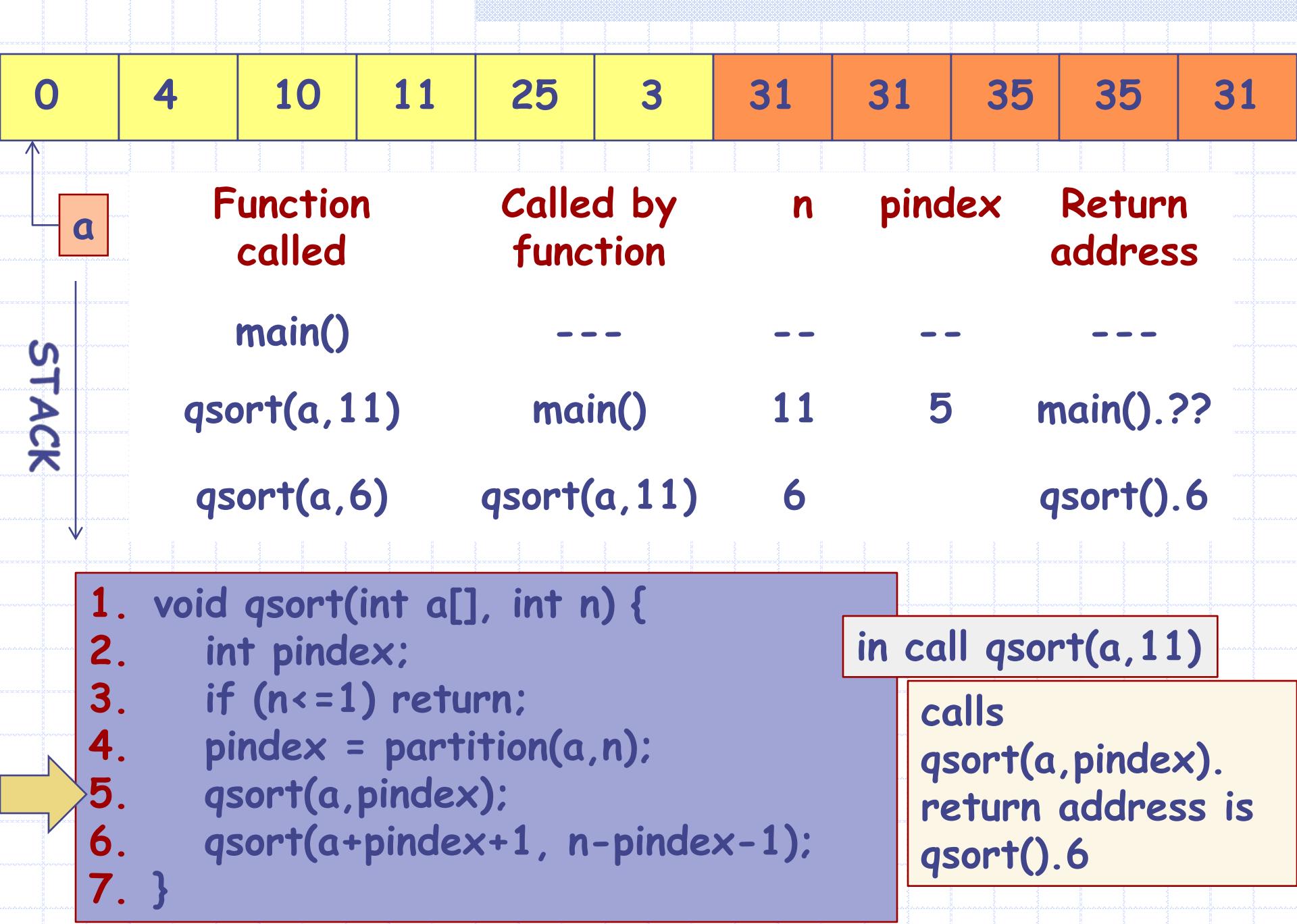
```
1. void qsort(int a[], int n) {  
2.     int pindex;  
3.     if (n<=1) return;  
4.     pindex = partition(a,n);  
5.     qsort(a,pindex);  
6.     qsort(a+pindex+1, n-pindex-1);  
7. }
```

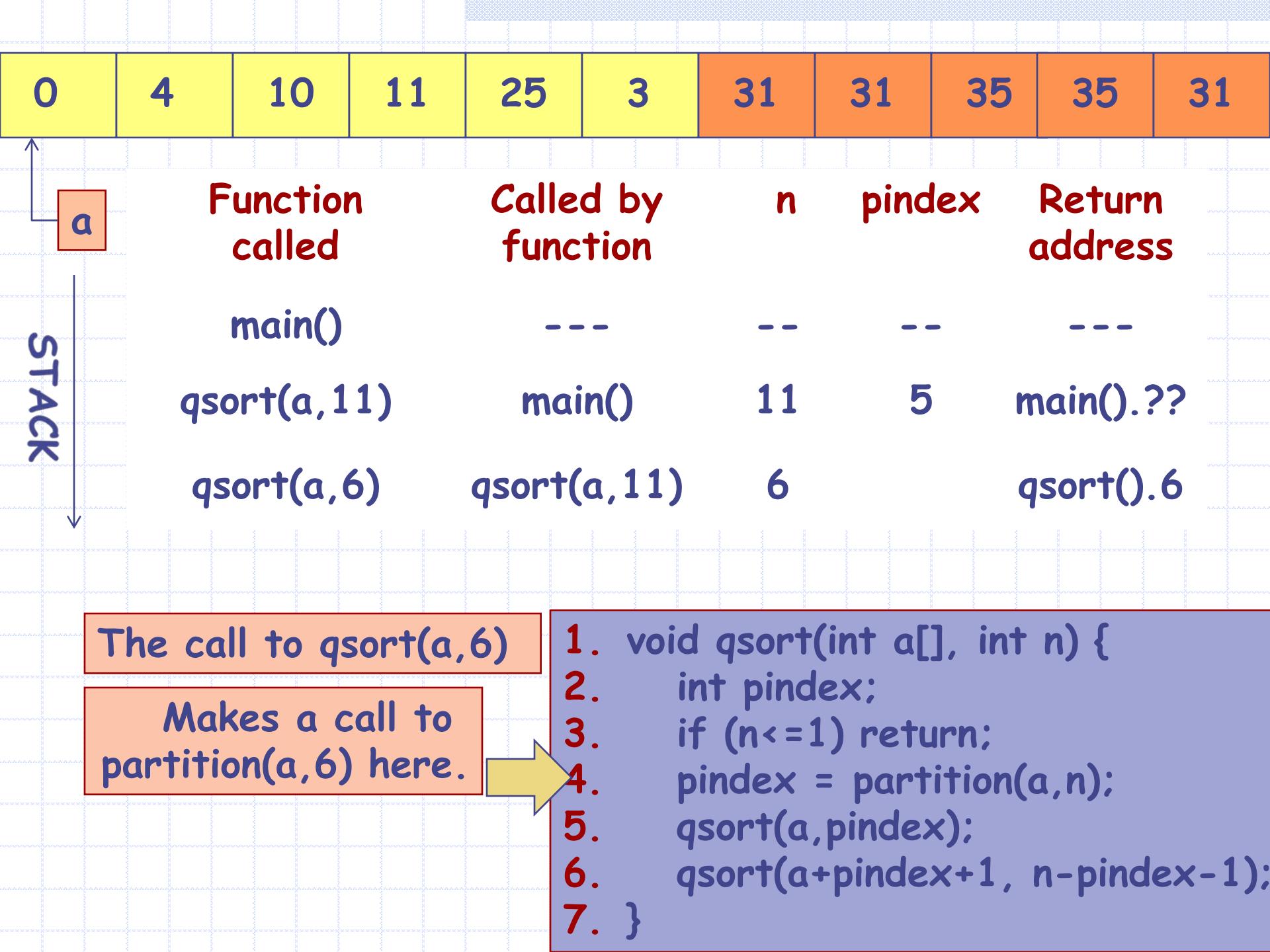


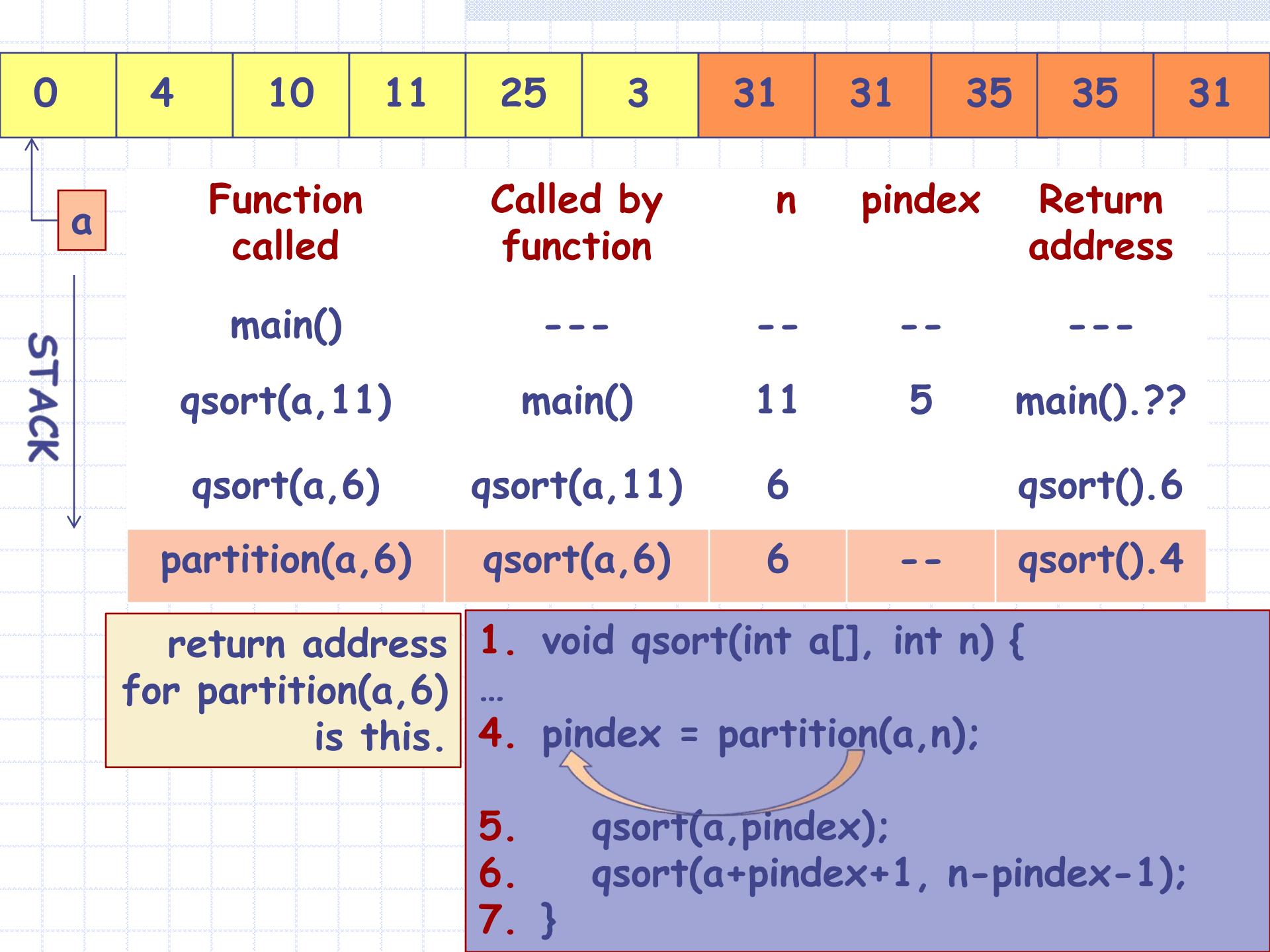
Quicksort function

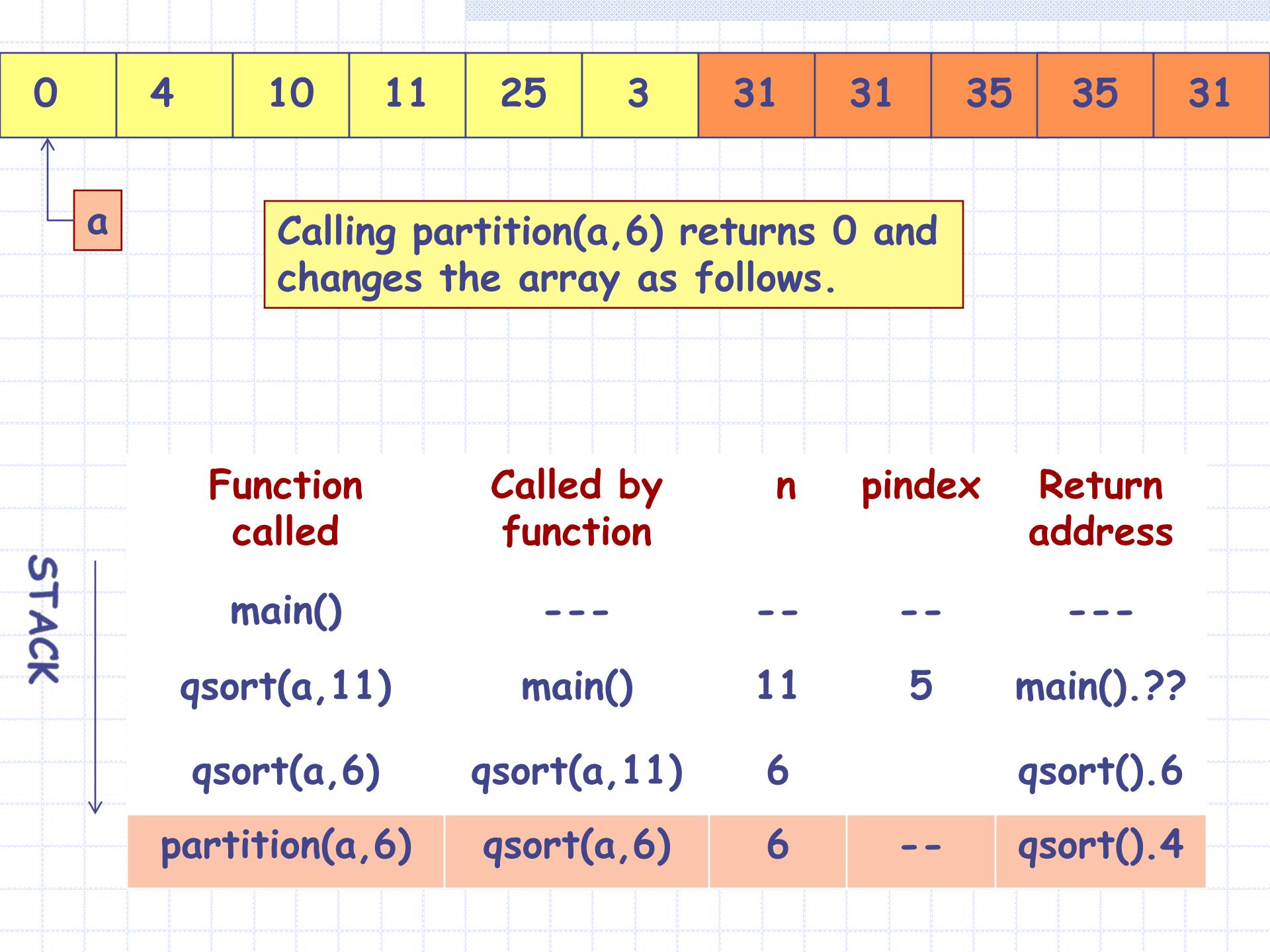
calls  
partition  
(a, 11)

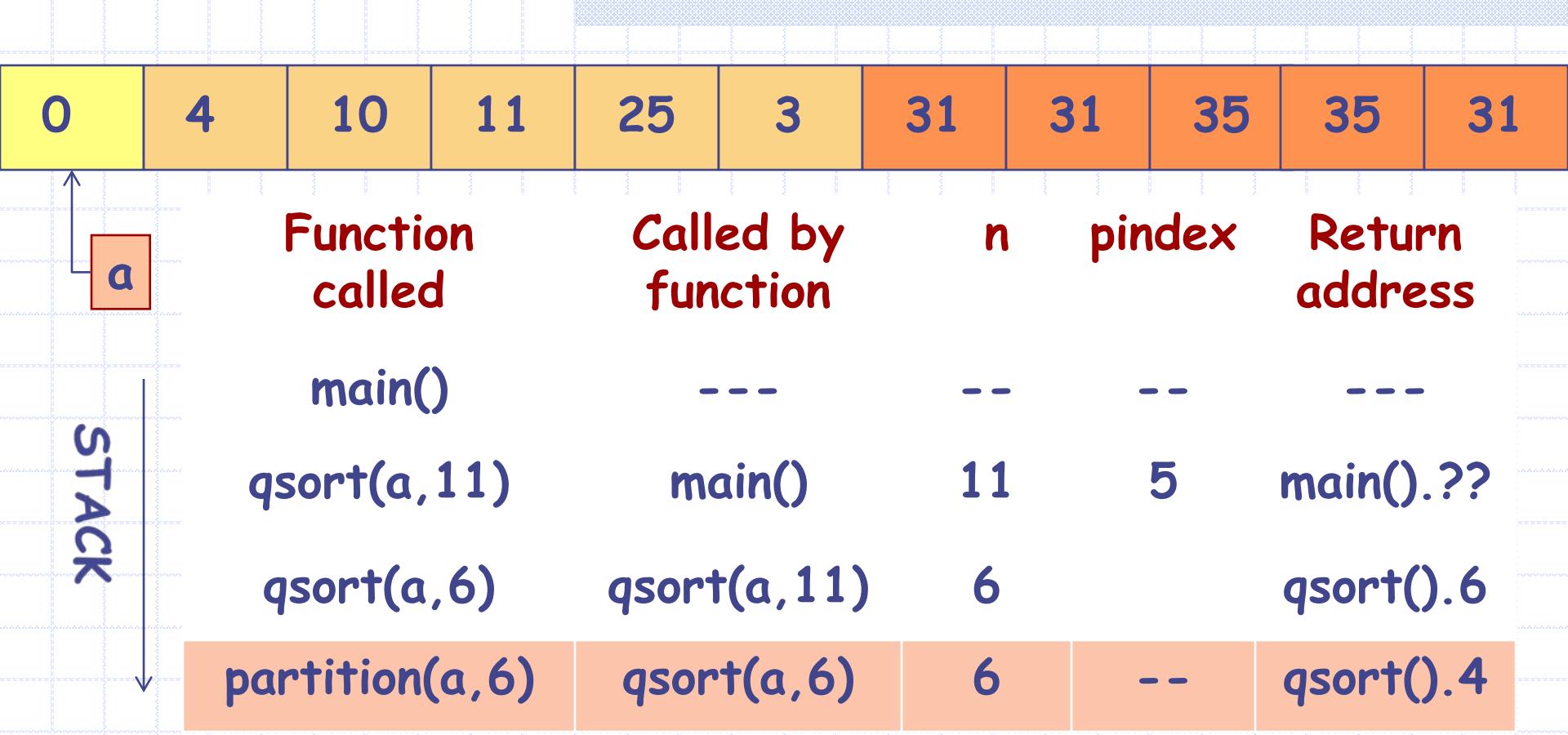












1. partition(a, 6)
2. returns.
3. Return value is 0.
4. pindex is set to 0.
5. qsort(a, 6)
6. resumes at line 4.

```

1. void qsort(int a[], int n) {
2.     int pindex;
3.     if (n<=1) return;
4.     pindex = partition(a,n);
5.     qsort(a,pindex);
6.     qsort(a+pindex+1, n-pindex-1);
7. }
```

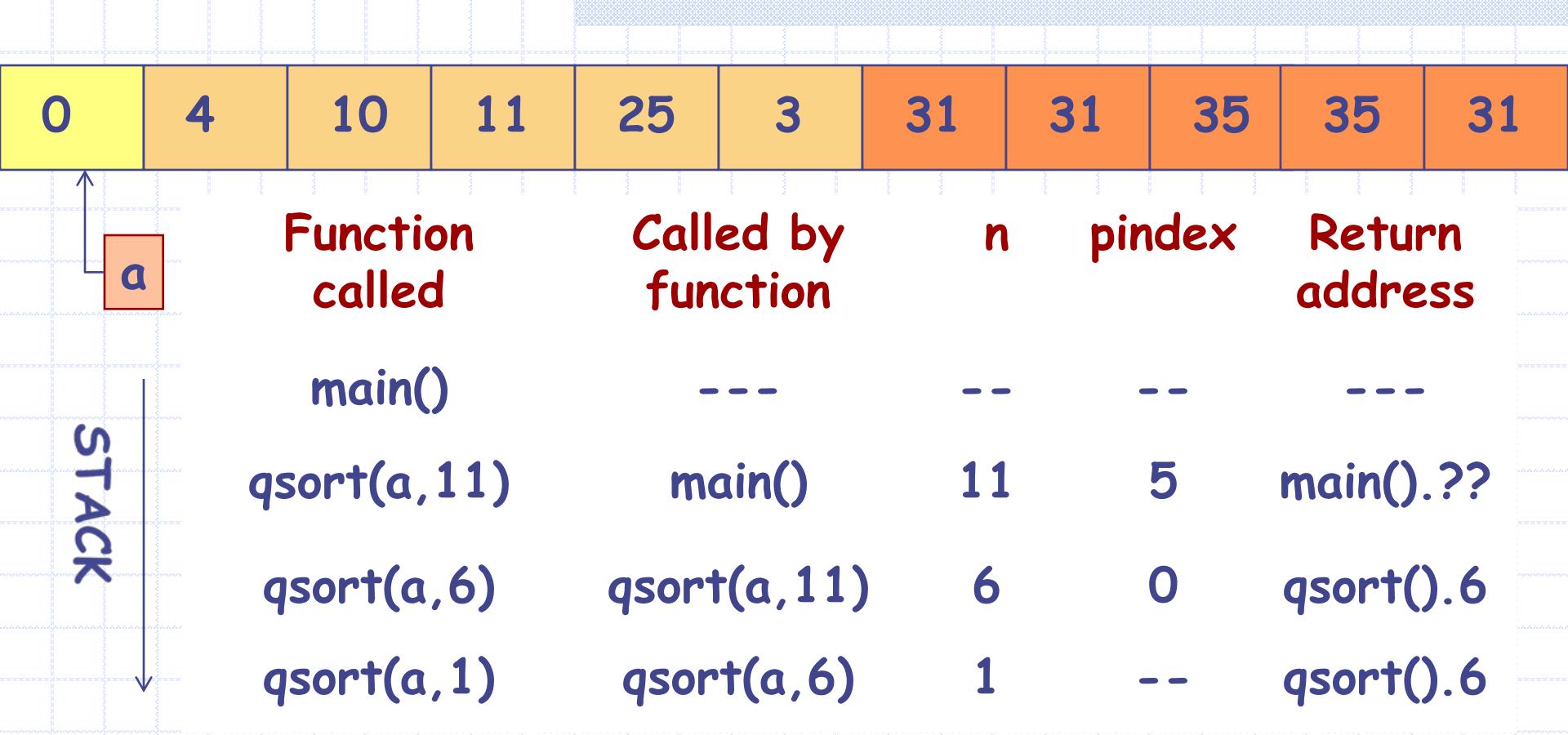
0	4	10	11	25	3	31	31	35	35	31	
a											
Function called					Called by function		n	pindex	Return address		
main()					---		--	--	---		
qsort(a, 11)					main()		11	5	main().??		
qsort(a, 6)					qsort(a, 11)		6	0	qsort().6		
qsort(a, 1)					qsort(a, 6)		1	--	qsort().6		

In call qsort(a, 6)

qsort(a, 6) now has pindex as 1.  
Now calls qsort(a, 1).  
Return addr. qsort()  
line 6.

```

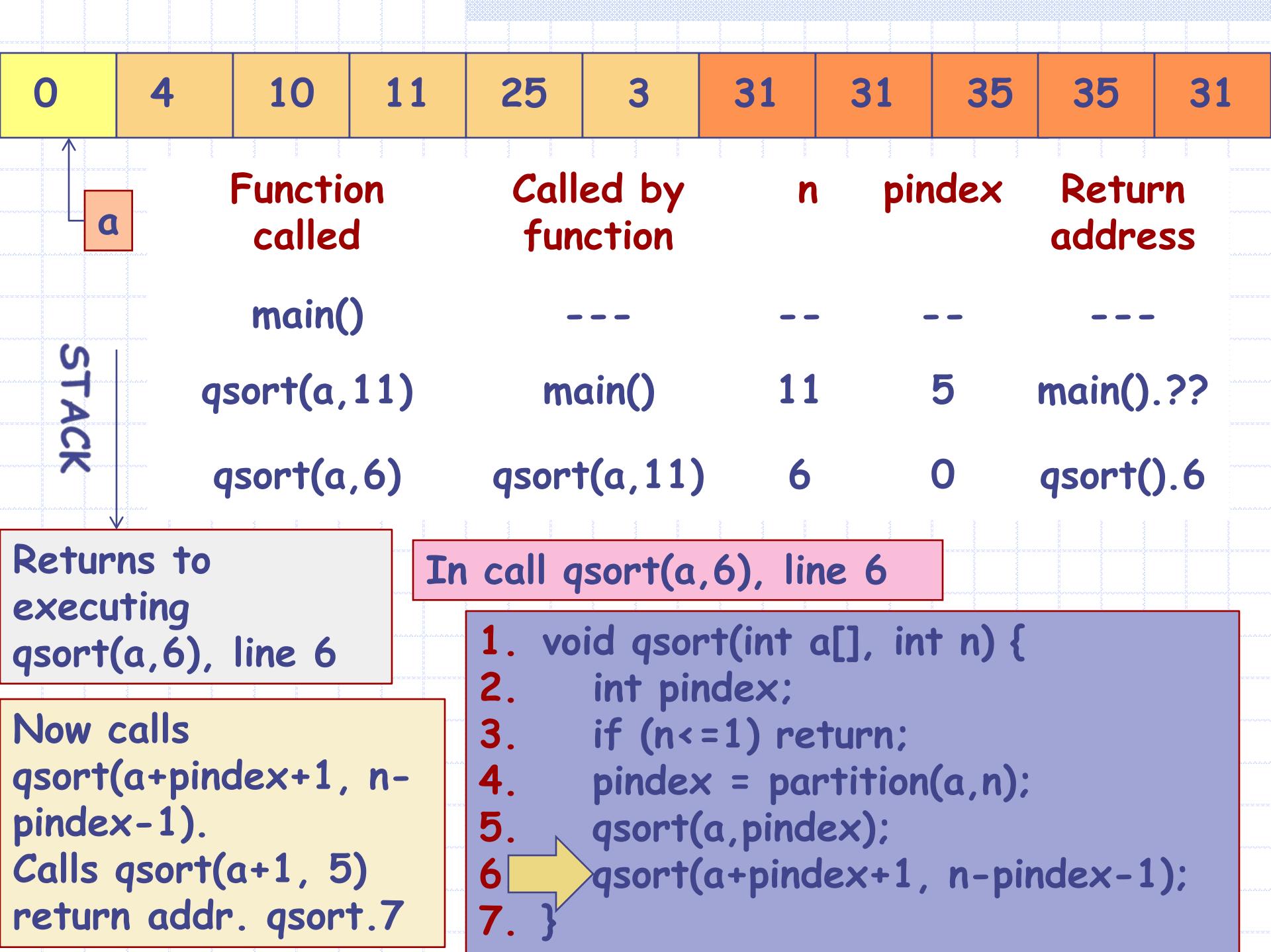
1. void qsort(int a[], int n) {
2.     int pindex;
3.     if (n<=1) return;
4.     pindex = partition(a,n);
5.     qsort(a,pindex);
6.     qsort(a+pindex+1, n-pindex-1);
7. }
```

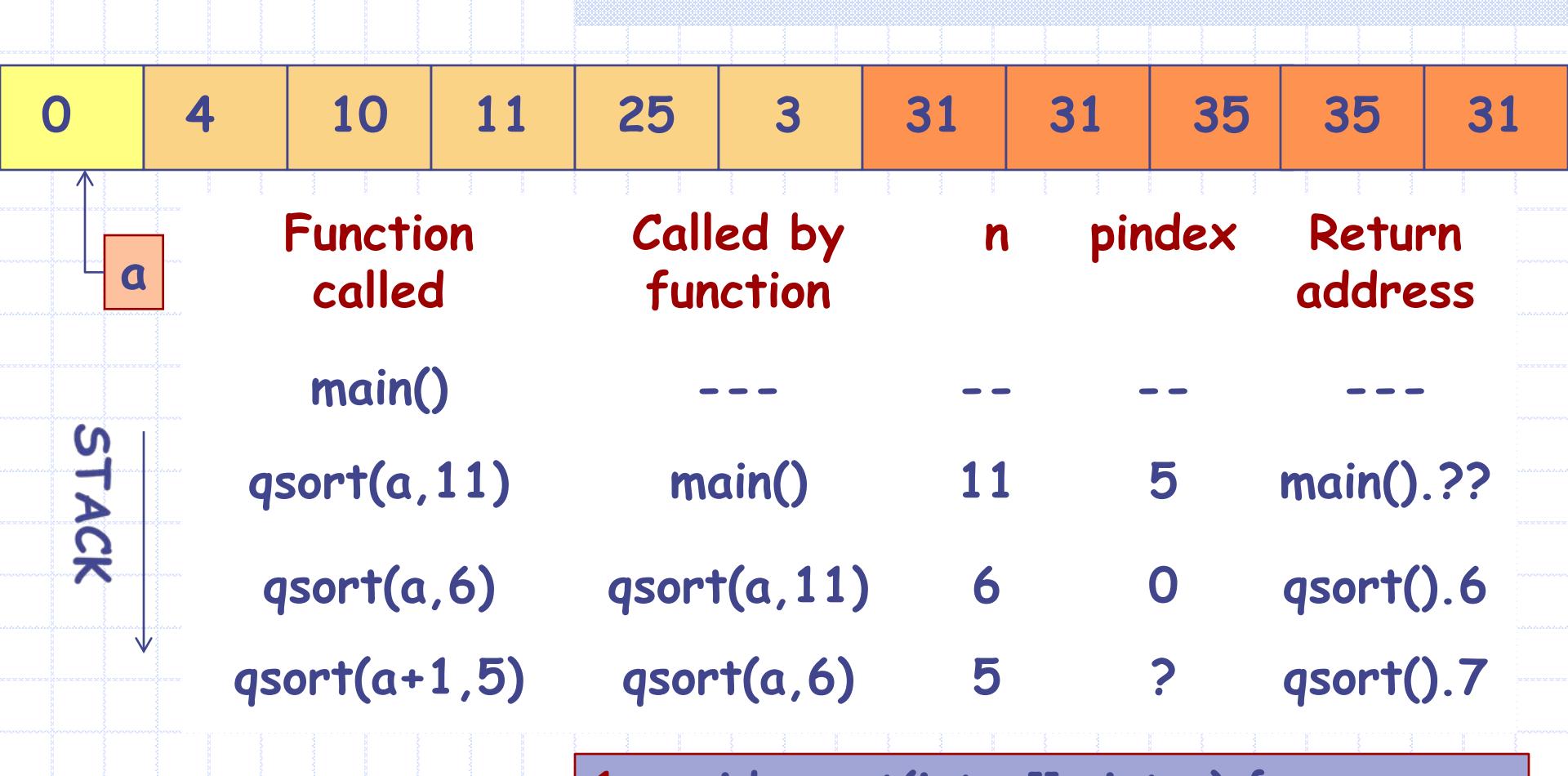


In call qsort(a, 1)

Since n is 1,  
qsort(a, 1) returns  
immediately.

1. void qsort(int a[], int n) {
2.     int pindex;
3.     if (n<=1) return;
4.     pindex = partition(a,n);
5.     qsort(a,pindex);
6.     qsort(a+pindex+1, n-pindex-1);
7. }





In call qsort(a+1, 5),  
line 3

Now calls  
partition(a+1, 5)

```

1. void qsort(int a[], int n) {
2.     int pindex;
3.     if (n<=1) return;
4.     pindex = partition(a,n);
5.     qsort(a,pindex);
6.     qsort(a+pindex+1, n-pindex-1);
7. }
```

	Function called	Called by function	n	pindex	Return address
	main()	---	--	--	---
qsort(a, 11)	main()	11	5	main().??	
qsort(a, 6)	qsort(a, 11)	6	0	qsort().6	
qsort(a+1, 5)	qsort(a, 6)	5	?	qsort().7	
partition(a+1, 5)	qsort(a+1, 5)	5	--	qsort().4	

In partition(a+1, 5):

Pivot is  $(a+1)[0]$  which is 4, so after partition is over, the array would be like this...

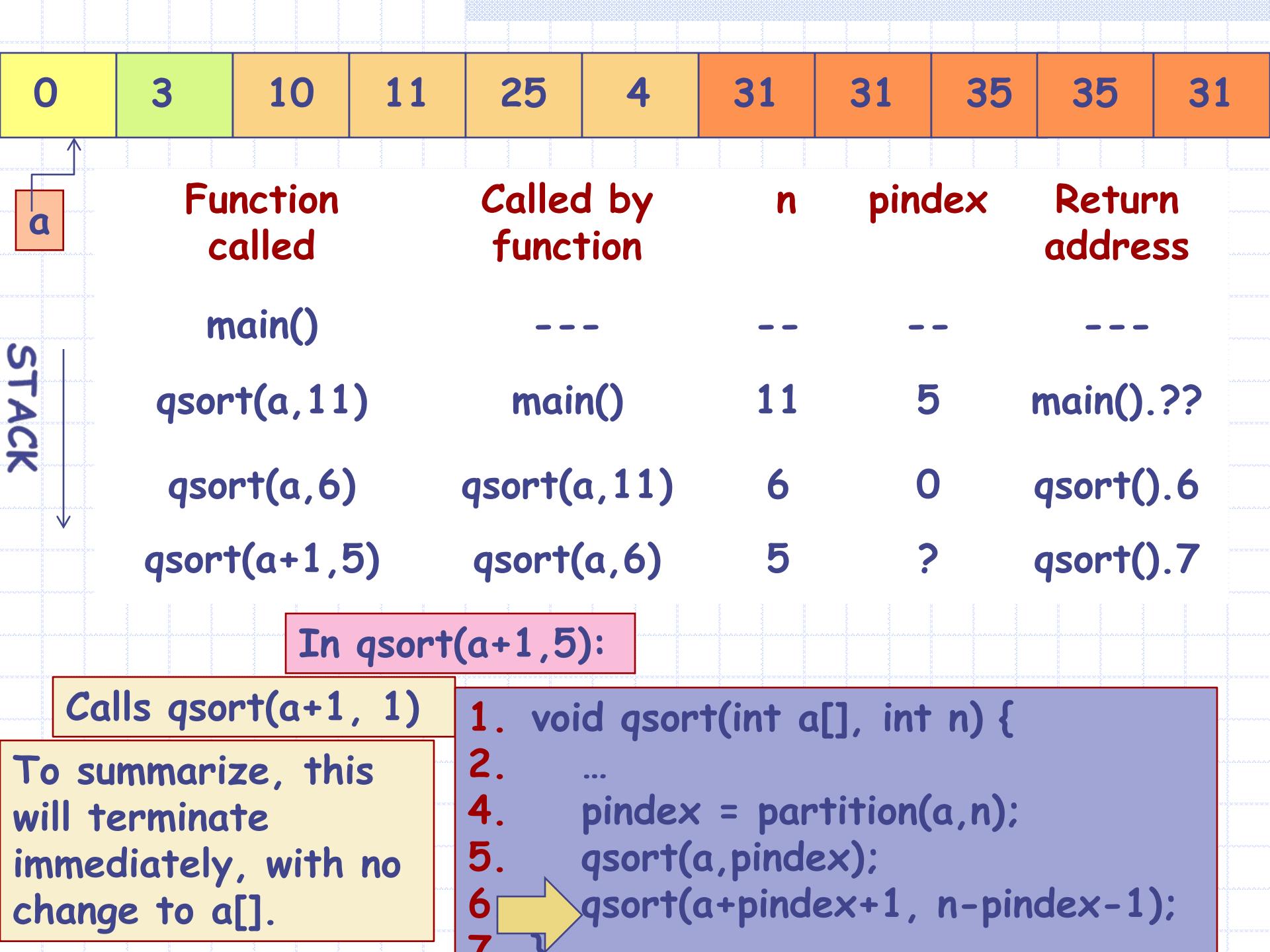
	0	3	10	11	25	4	31	31	35	35	31
Function called	a	Called by function		n	pindex	Return address					
main()		---	--	--	--	---					
qsort(a, 11)		main()	11	5	main().??						
qsort(a, 6)		qsort(a, 11)	6	0	qsort().6						
qsort(a+1, 5)		qsort(a, 6)	5	?	qsort().7						

In qsort(a+1, 5):

partition(a+1, 5)  
returns 0. So  
pindex is 0.  
execution resumes  
at line 4.

```

1. void qsort(int a[], int n) {
2. ...
4.     pindex = partition(a,n);
5.     qsort(a,pindex);
6.     qsort(a+pindex+1, n-pindex-1);
7. }
```



0	3	10	11	25	4	31	31	35	35	31
a										
Function called										
main()	---									---
qsort(a, 11)	main()									main().??
qsort(a, 6)	qsort(a, 11)									qsort().6
qsort(a+1, 5)	qsort(a, 6)									qsort().6

In qsort(a+1, 5):

in line 6: pindex is 0  
Calls qsort(a+2, 4).

```

1. void qsort(int a[], int n) {
2. ...
4.     pindex = partition(a,n);
5.     qsort(a,pindex);
6.     qsort(a+pindex+1, n-pindex-1);
7.

```

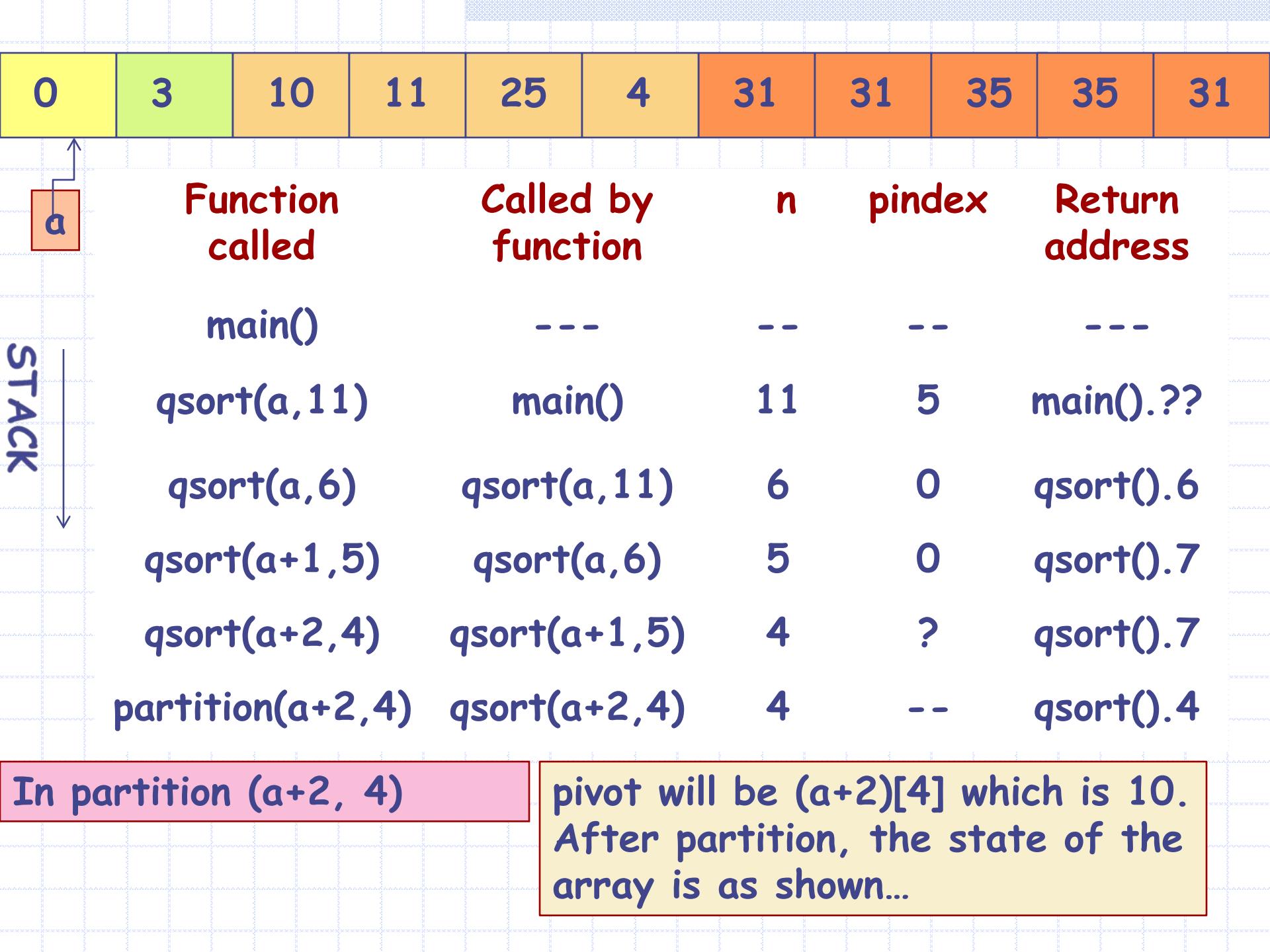
	0	3	10	11	25	4	31	31	35	35	31
Function called	a										
n											Return address
main()	---										---
qsort(a, 11)		main()					11	5	main().??		
qsort(a, 6)		qsort(a, 11)				6	0	0	qsort().6		
qsort(a+1, 5)		qsort(a, 6)				5	0	0	qsort().7		
qsort(a+2, 4)		qsort(a+1, 5)				4	?	?	qsort().7		

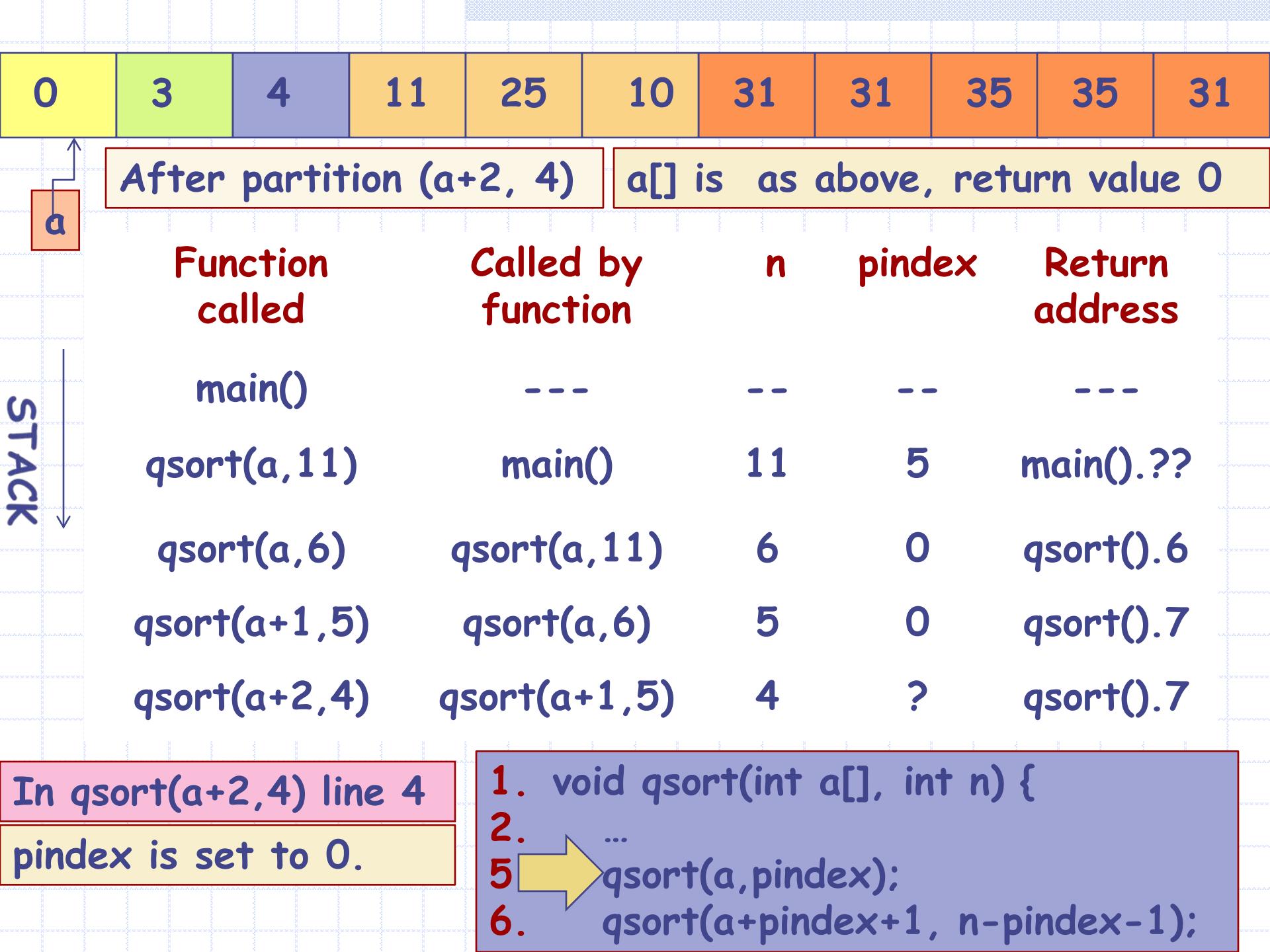
in line 6: pindex is 0  
Calls qsort(a+2, 4).  
return addr is  
qsort().line 7

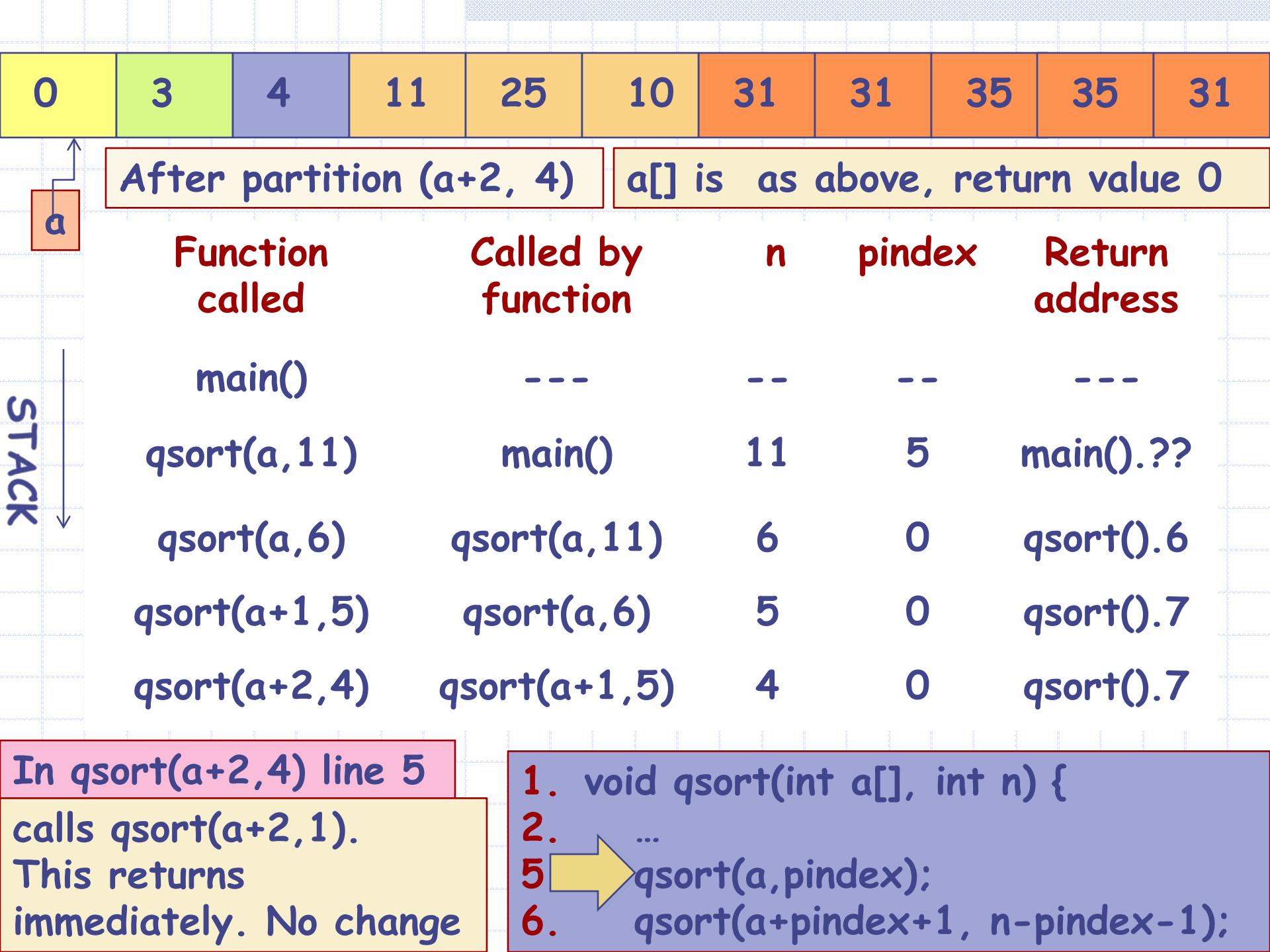
```

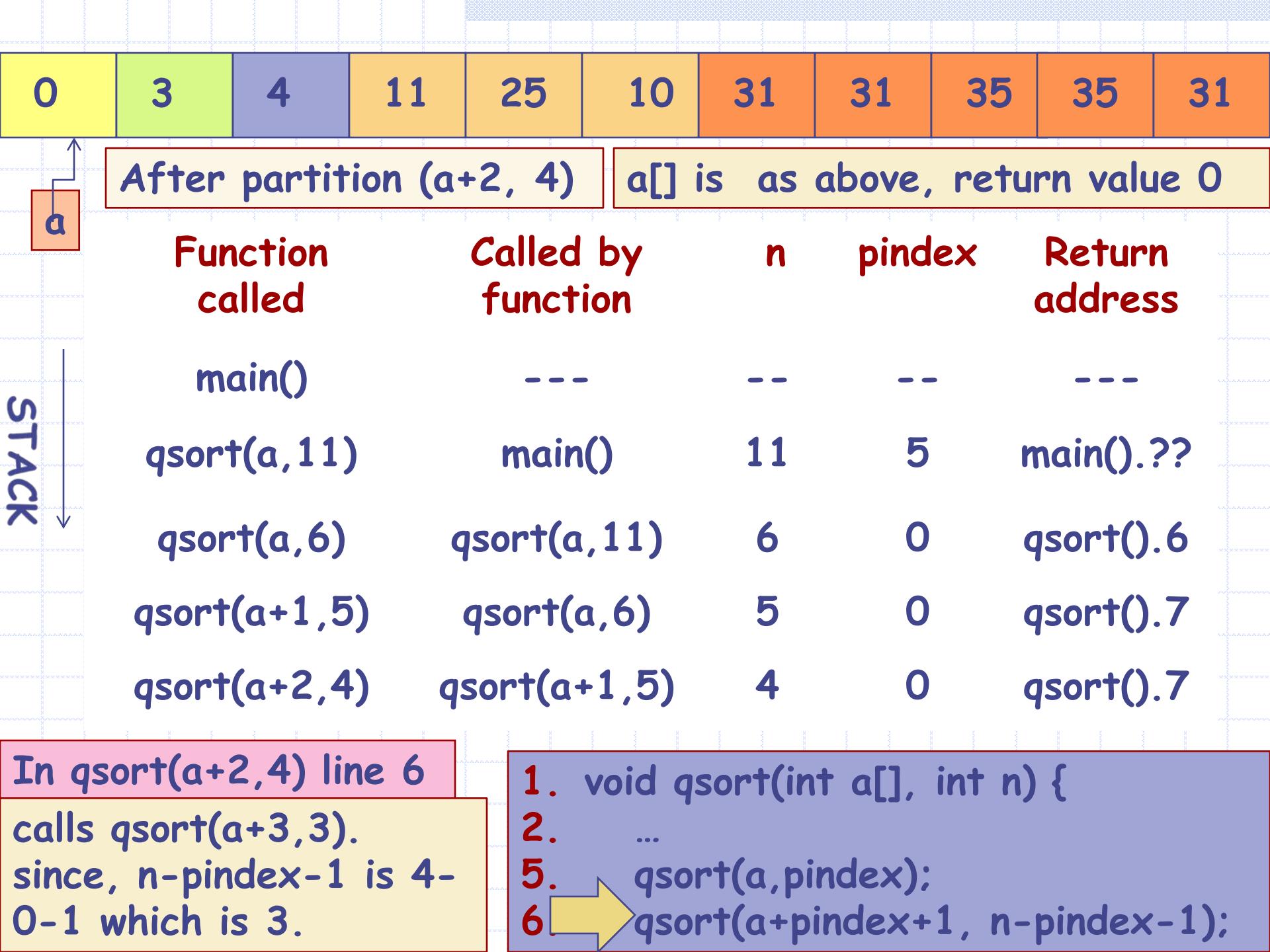
1. void qsort(int a[], int n) {
2. ...
4.     pindex = partition(a,n);
5.     qsort(a,pindex);
6.     qsort(a+pindex+1, n-pindex-1);
7.

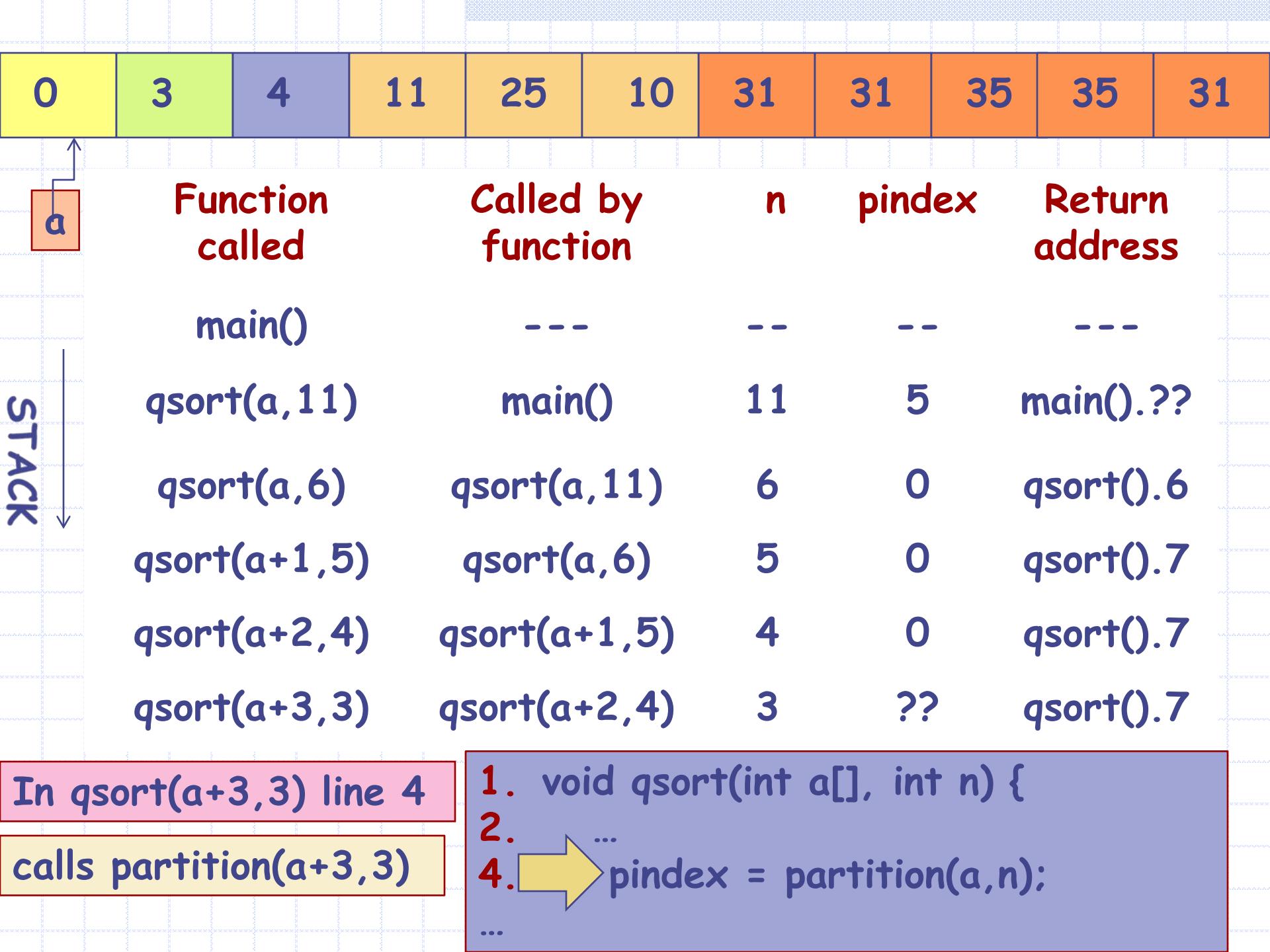
```









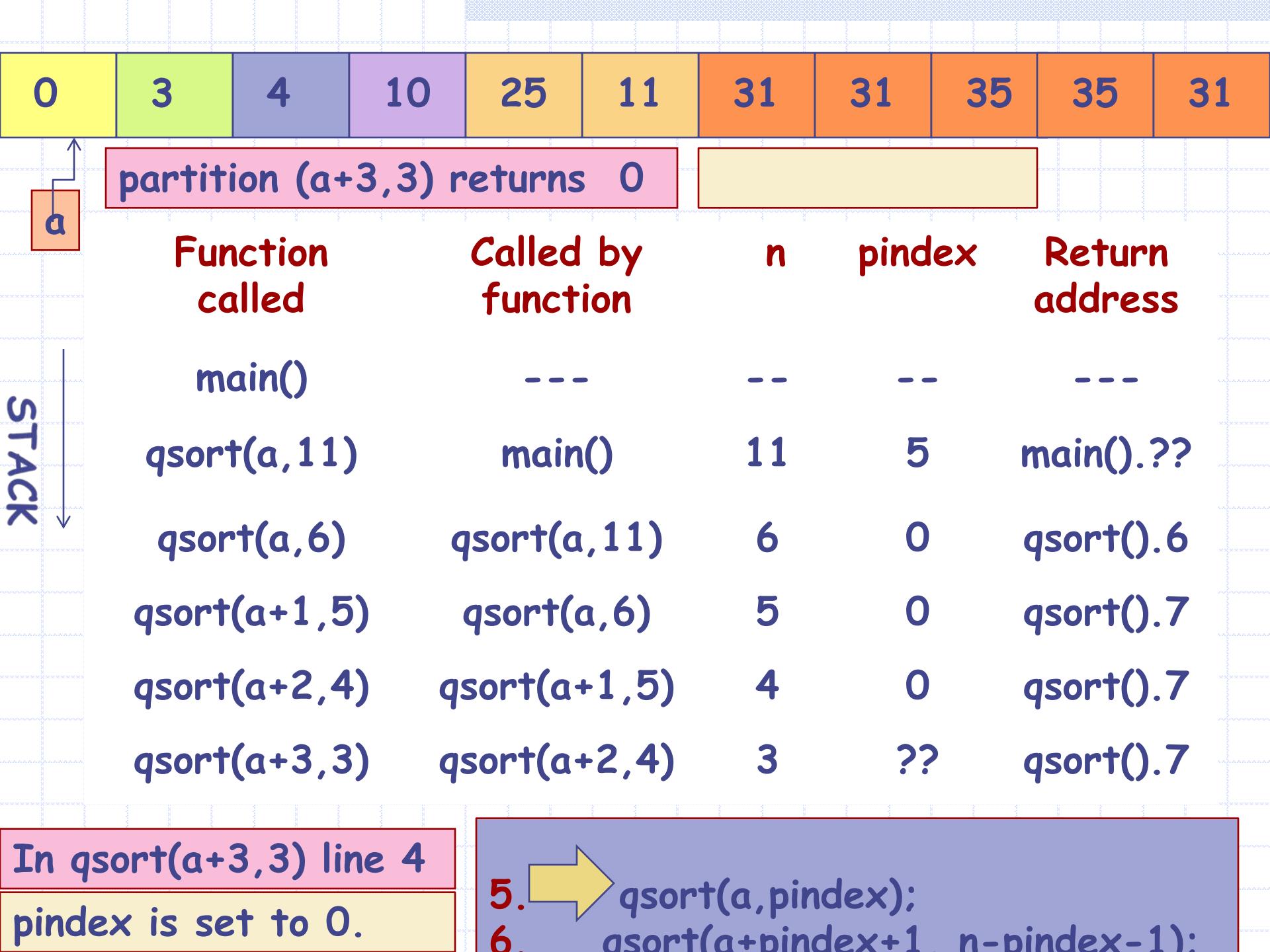


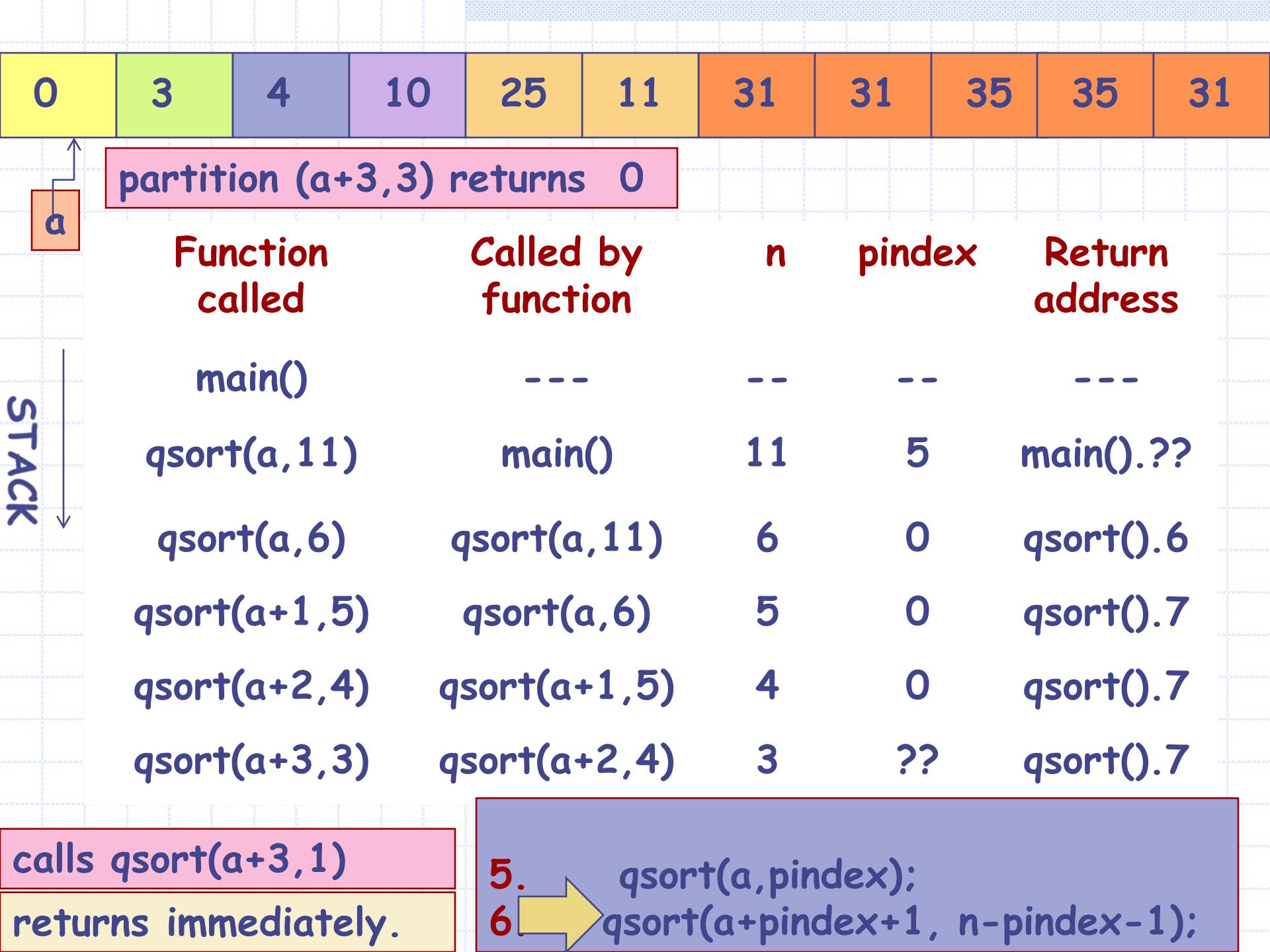
0	3	4	11	25	10	31	31	35	35	31
a	Function called	Called by function	n	pindex	Return address					
main()		---	--	--	---					
qsort(a, 11)	main()	11	5	main().??						
qsort(a, 6)	qsort(a, 11)	6	0	qsort().6						
qsort(a+1, 5)	qsort(a, 6)	5	0	qsort().7						
qsort(a+2, 4)	qsort(a+1, 5)	4	0	qsort().7						
qsort(a+3, 3)	qsort(a+2, 4)	3	??	qsort().7						
partition(a+3, 3)	qsort(a+3, 3)	3	--	qsort().4						

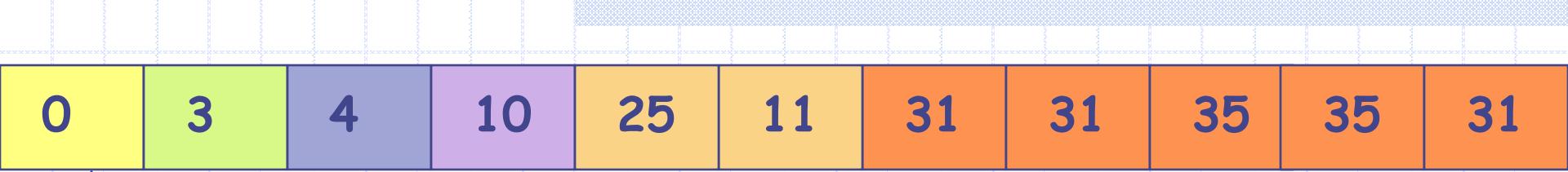
In partition  
(a+3, 3)

pivot is (a+3)[0]  
which is 11.

State of array after  
partition becomes







a

partition ( $a+3, 3$ ) returns 0

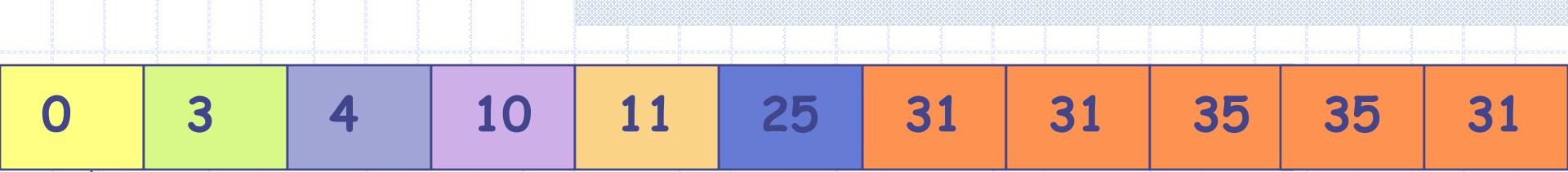
pindex is set to 0

Function called	Called by function	n	pindex	Return address
main()	---	--	--	---
qsort(a, 11)	main()	11	5	main().??
qsort(a, 6)	qsort(a, 11)	6	0	qsort().6
qsort(a+1, 5)	qsort(a, 6)	5	0	qsort().7
qsort(a+2, 4)	qsort(a+1, 5)	4	0	qsort().7
qsort(a+3, 3)	qsort(a+2, 4)	3	0	qsort().7

calls qsort( $a+4, 2$ )

5.  $\rightarrow$  qsort( $a, pindex$ );  
 6.  $\rightarrow$  qsort( $a+pindex+1, n-pindex-1$ );

0	3	4	10	25	11	31	31	35	35	31	
<b>a</b>	Function called	Called by function	n	pindex	Return address						
main()	---	---	--	--	---	---	---	---	---	---	---
qsort(a, 11)	main()	main()	11	5	main().??						
qsort(a, 6)	qsort(a, 11)	qsort(a, 11)	6	0	qsort().6						
qsort(a+1, 5)	qsort(a, 6)	qsort(a, 6)	5	0	qsort().7						
qsort(a+2, 4)	qsort(a+1, 5)	qsort(a+1, 5)	4	0	qsort().7						
qsort(a+3, 3)	qsort(a+2, 4)	qsort(a+2, 4)	3	0	qsort().7						
qsort(a+4, 2)	qsort(a+3, 3)	qsort(a+3, 3)	2	??	qsort().7						
partition(a+4, 2)	qsort(a+4, 2)	qsort(a+4, 2)	2	--	qsort().4						
calls partition (a+4, 2)				pivot is 25				array becomes...			



a

partition (a+4, 2) returns 0

pindex is 0

Function called

Called by function

n

pindex

Return address

main()

---

--

--

---

STACK

qsort(a, 11)

main()

11

5

main().??

qsort(a, 6)

qsort(a, 11)

6

0

qsort().6

qsort(a+1, 5)

qsort(a, 6)

5

0

qsort().7

qsort(a+2, 4)

qsort(a+1, 5)

4

0

qsort().7

qsort(a+3, 3)

qsort(a+2, 4)

3

0

qsort().7

qsort(a+4, 2)

qsort(a+3, 3)

2

??

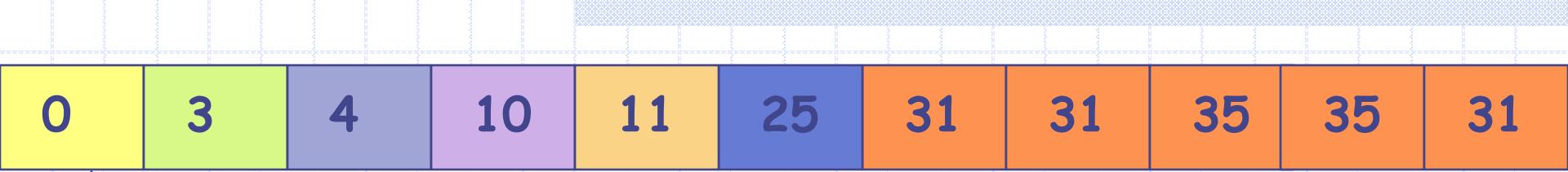
qsort().7

qsort(a+4, 2) resumes at line 7. But this is the last line...

```
1. void qsort(int a[], int n) {  
2.     int pindex;  
3.     if (n<=1) return;  
4.     pindex = partition(a,n);  
5.     qsort(a,pindex);  
6.     qsort(a+pindex+1, n-pindex-1);  
7. }
```

Line 7 terminates the call to `qsort(a, n)`.

So stack changes as follows.



a

partition ( $a+4, 2$ ) returns 0

pindex is 0

Function called	Called by function	n	pindex	Return address
-----------------	--------------------	---	--------	----------------

main()

---

n

pindex

Return address

qsort(a, 11)

main()

--

--

main().??

qsort(a, 6)

qsort(a, 11)

11

5

qsort().6

qsort(a+1, 5)

qsort(a, 6)

6

0

qsort().7

qsort(a+2, 4)

qsort(a+1, 5)

5

0

qsort().7

qsort(a+3, 3)

qsort(a+2, 4)

4

0

qsort().7

qsort(a+4, 2)

qsort(a+3, 3)

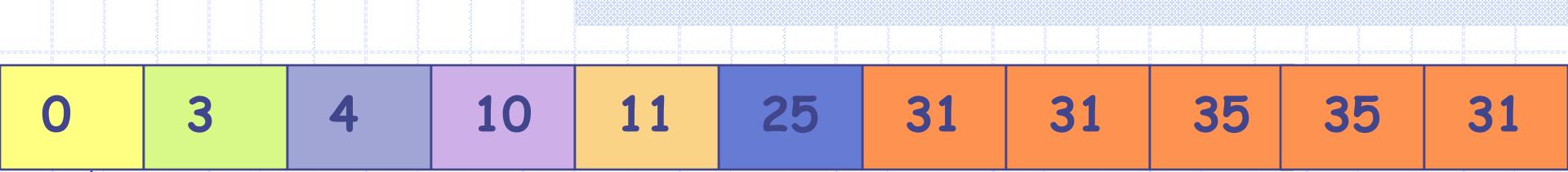
3

0

qsort().7

qsort(a+4, 2) resumes at line 4, pindex is 0.

STACK



<b>a</b>	<b>Function called</b>	<b>Called by function</b>	<b>n</b>	<b>pindex</b>	<b>Return address</b>
	main()		---	--	---
	qsort(a, 11)	main()	11	5	main().??
STACK	qsort(a, 6)	qsort(a, 11)	6	0	qsort().6
	qsort(a+1, 5)	qsort(a, 6)	5	0	qsort().7
	qsort(a+2, 4)	qsort(a+1, 5)	4	0	qsort().7
	qsort(a+3, 3)	qsort(a+2, 4)	3	0	qsort().7
	qsort(a+4, 2)	qsort(a+3, 3)	2	0	qsort().7

**qsort(a+4, 2) line 5: Calls qsort(a+5, 1) which terminates immediately.**

Function called	Called by function	n	pindex	Return address
main()	---	--	--	---
qsort(a, 11)	main()	11	5	main().??
qsort(a, 6)	qsort(a, 11)	6	0	qsort().6
qsort(a+1, 5)	qsort(a, 6)	5	0	qsort().7
qsort(a+2, 4)	qsort(a+1, 5)	4	0	qsort().7
qsort(a+3, 3)	qsort(a+2, 4)	3	0	qsort().7
qsort(a+4, 2)	qsort(a+3, 3)	2	0	qsort().7

qsort(a+4, 2) line 6: Calls qsort(a+6, 1) which terminates immediately.

Function called	Called by function	n	pindex	Return address
main()	---	--	--	---
qsort(a, 11)	main()	11	5	main().??
qsort(a, 6)	qsort(a, 11)	6	0	qsort().6
qsort(a+1, 5)	qsort(a, 6)	5	0	qsort().7
qsort(a+2, 4)	qsort(a+1, 5)	4	0	qsort().7
qsort(a+3, 3)	qsort(a+2, 4)	3	0	qsort().7
qsort(a+4, 2)	qsort(a+3, 3)	2	0	qsort().7

qsort(a+4, 2) line 7: qsort(a+4, 2) terminates now.

Control returns to its calling fn: qsort(a+3, 3) at line 7.

Function called	Called by function	n	pindex	Return address
main()	---	--	--	---
qsort(a, 11)	main()	11	5	main().??
qsort(a, 6)	qsort(a, 11)	6	0	qsort().6
qsort(a+1, 5)	qsort(a, 6)	5	0	qsort().7
qsort(a+2, 4)	qsort(a+1, 5)	4	0	qsort().7
qsort(a+3, 3)	qsort(a+2, 4)	3	0	qsort().7

qsort(a+3, 3) resumes at line 7 and terminates. Control returns to calling fn qsort(a+2, 4) at line 7.

0	3	4	10	11	25	31	31	35	35	31
a										
Function called				Called by function		n	pindex	Return address		
main()				---		--	--	---		
qsort(a, 11)			main()			11	5	main().??		
qsort(a, 6)		qsort(a, 11)				6	0	qsort().6		
qsort(a+1, 5)	qsort(a, 6)					5	0	qsort().7		
qsort(a+2, 4)	qsort(a+1, 5)					4	0	qsort().7		

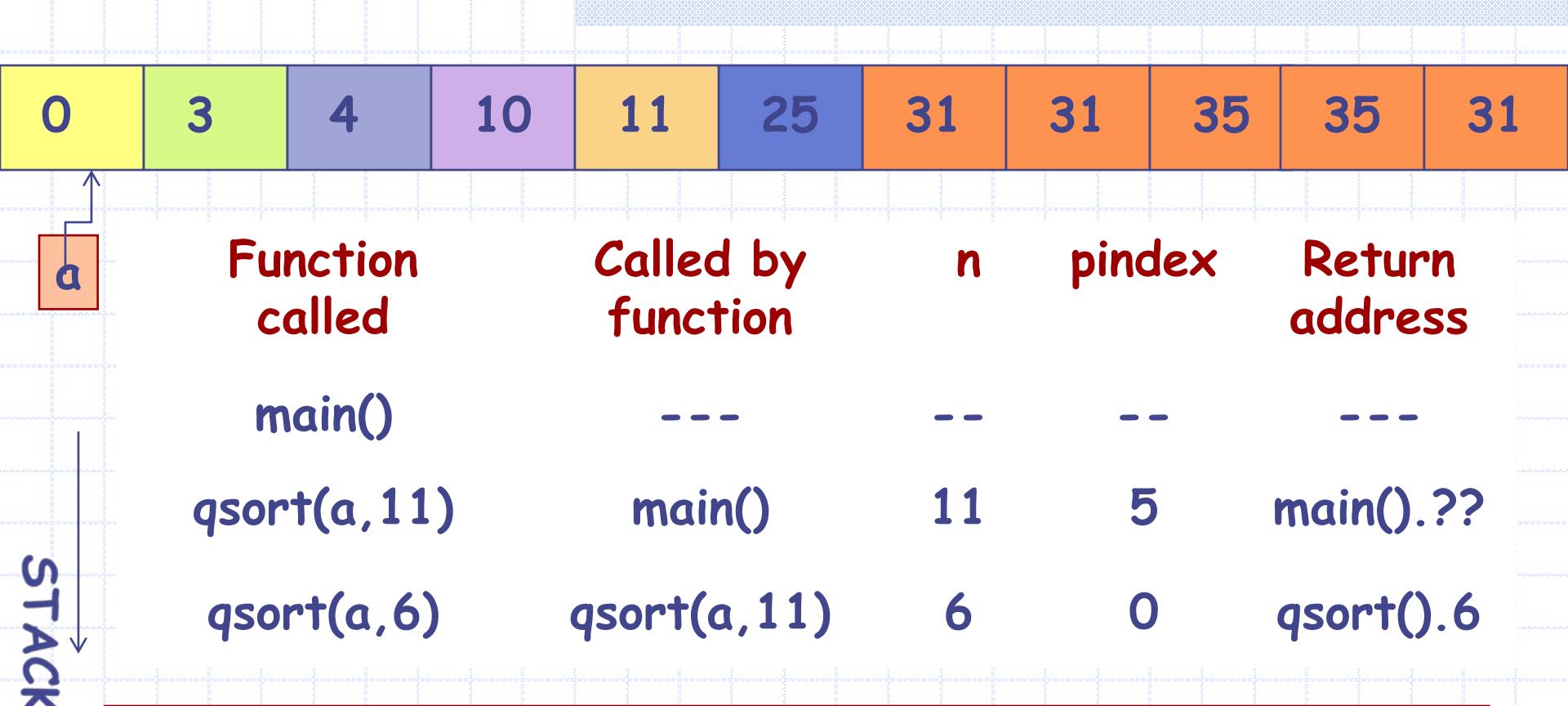
qsort(a+2, 4) resumes at line 7 and terminates. Control returns to calling fn qsort(a+1, 5) at line 7.

0	3	4	10	11	25	31	31	35	35	31
a	Function called	Called by function	n	pindex	Return address					
main()		---	--	--	---					
qsort(a, 11)	main()		11	5	main().??					
qsort(a, 6)	qsort(a, 11)		6	0	qsort().6					
qsort(a+1, 5)	qsort(a, 6)		5	0	qsort().7					

qsort(a+1, 5) resumes at line 7 and terminates. Control returns to calling fn qsort(a, 6) at line 7.

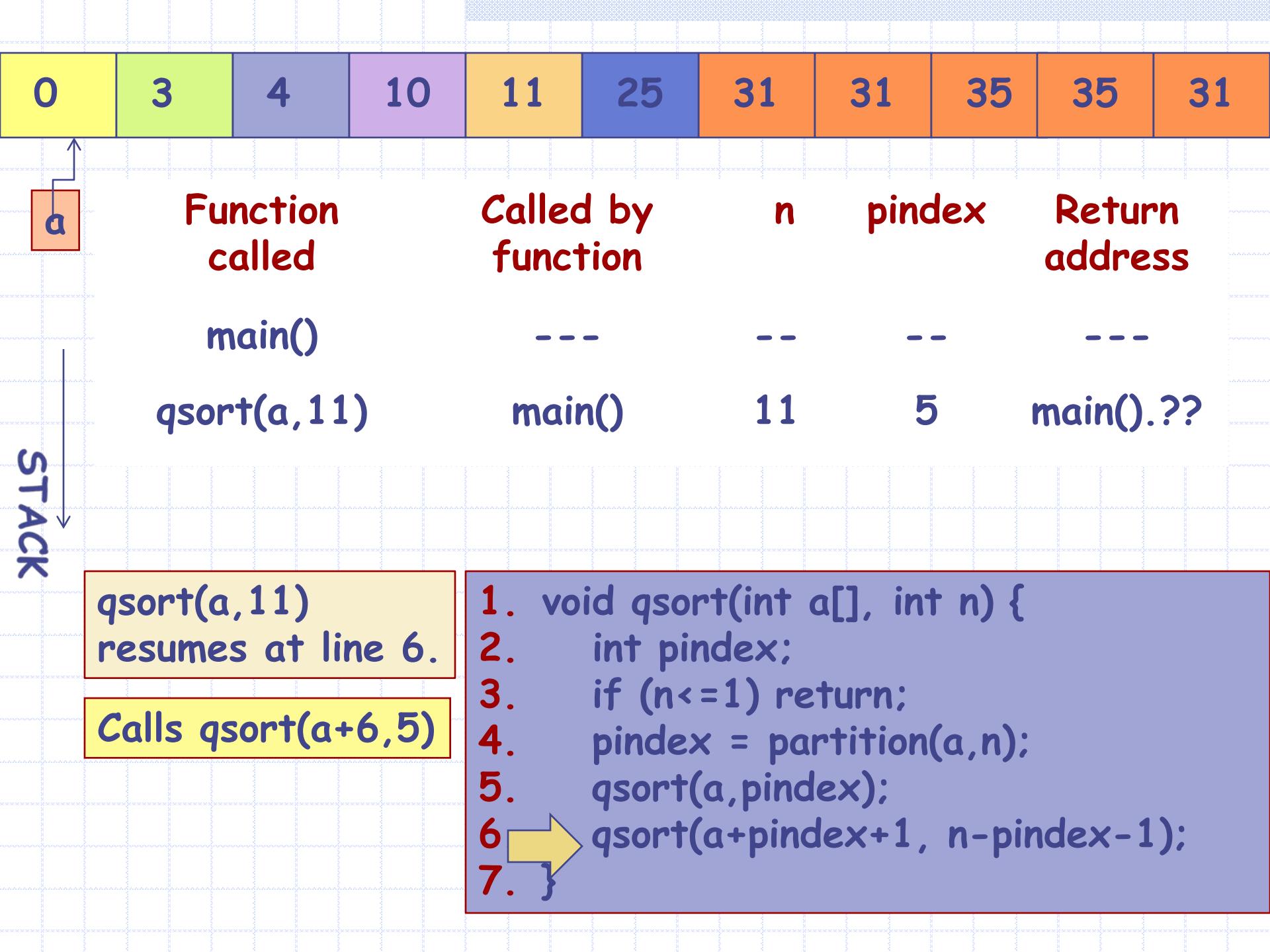
0	3	4	10	11	25	31	31	35	35	31
a	Function called	Called by function	n	pindex	Return address					
main()		---	--	--	---					
qsort(a, 11)	main()		11	5	main().??					
qsort(a, 6)	qsort(a, 11)		6	0	qsort().6					
qsort(a+1, 5)	qsort(a, 6)		5	0	qsort().7					

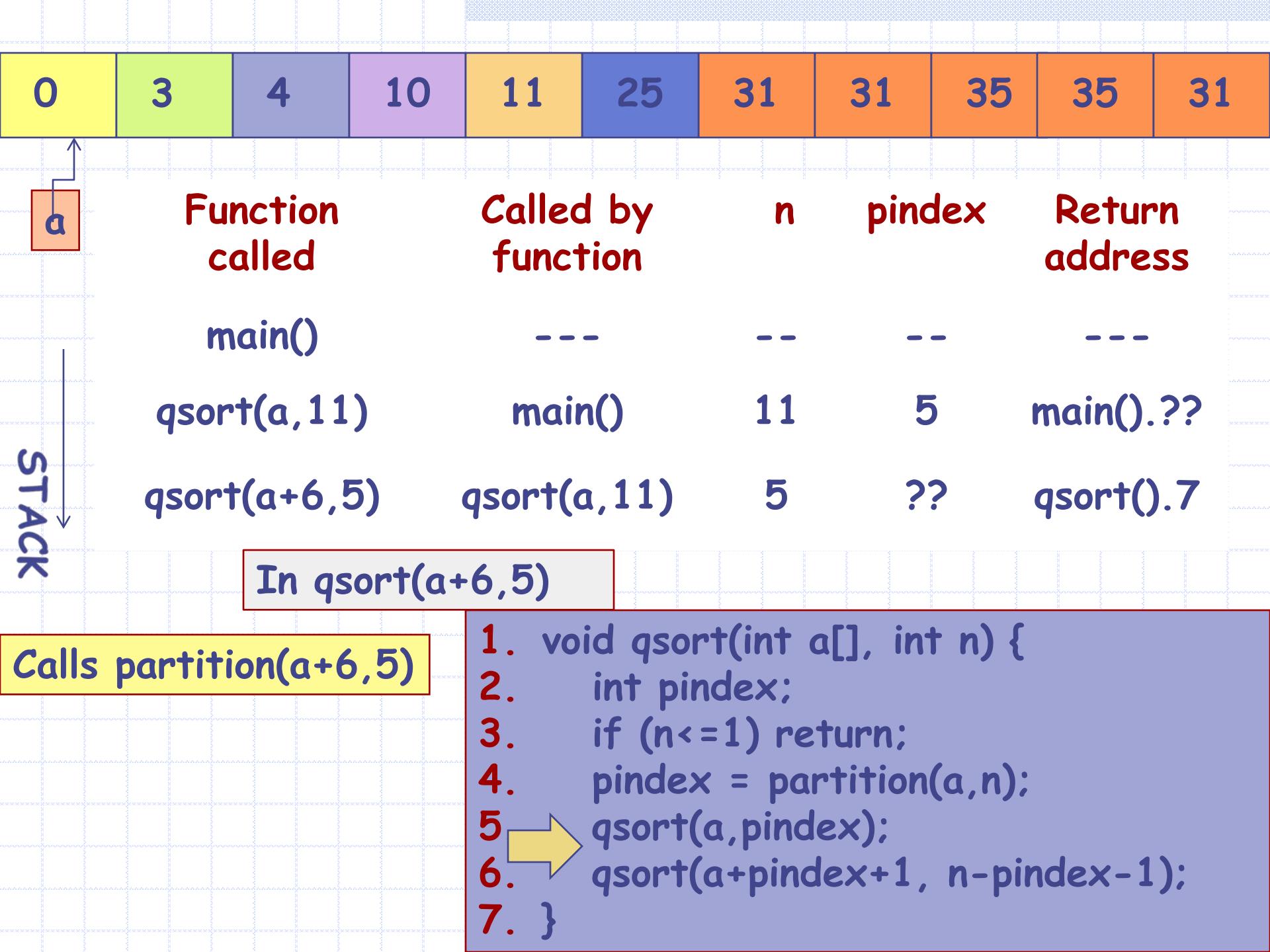
qsort(a+1, 5) resumes at line 7 and terminates. Control returns to calling fn qsort(a, 6) at line 7.

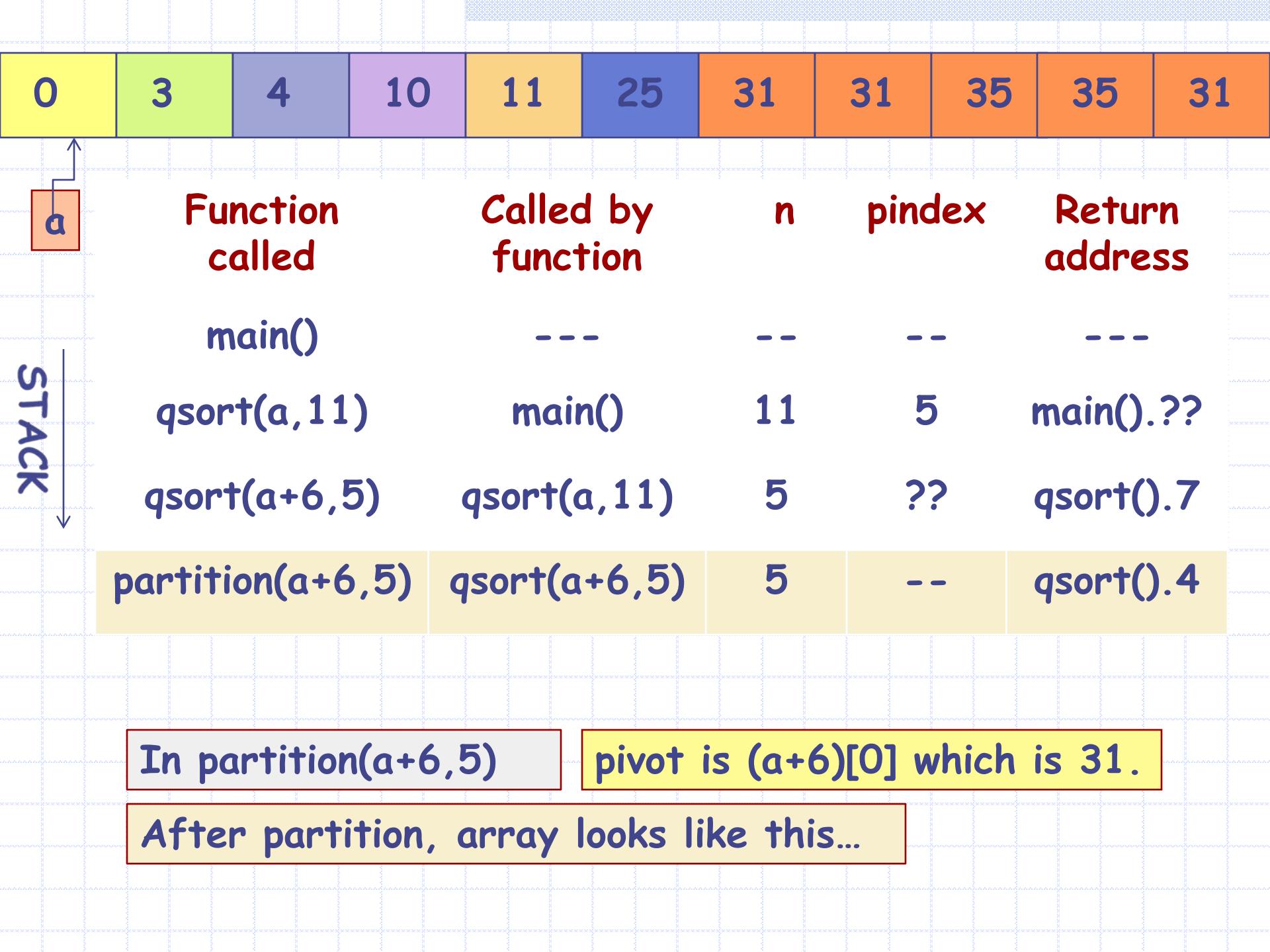


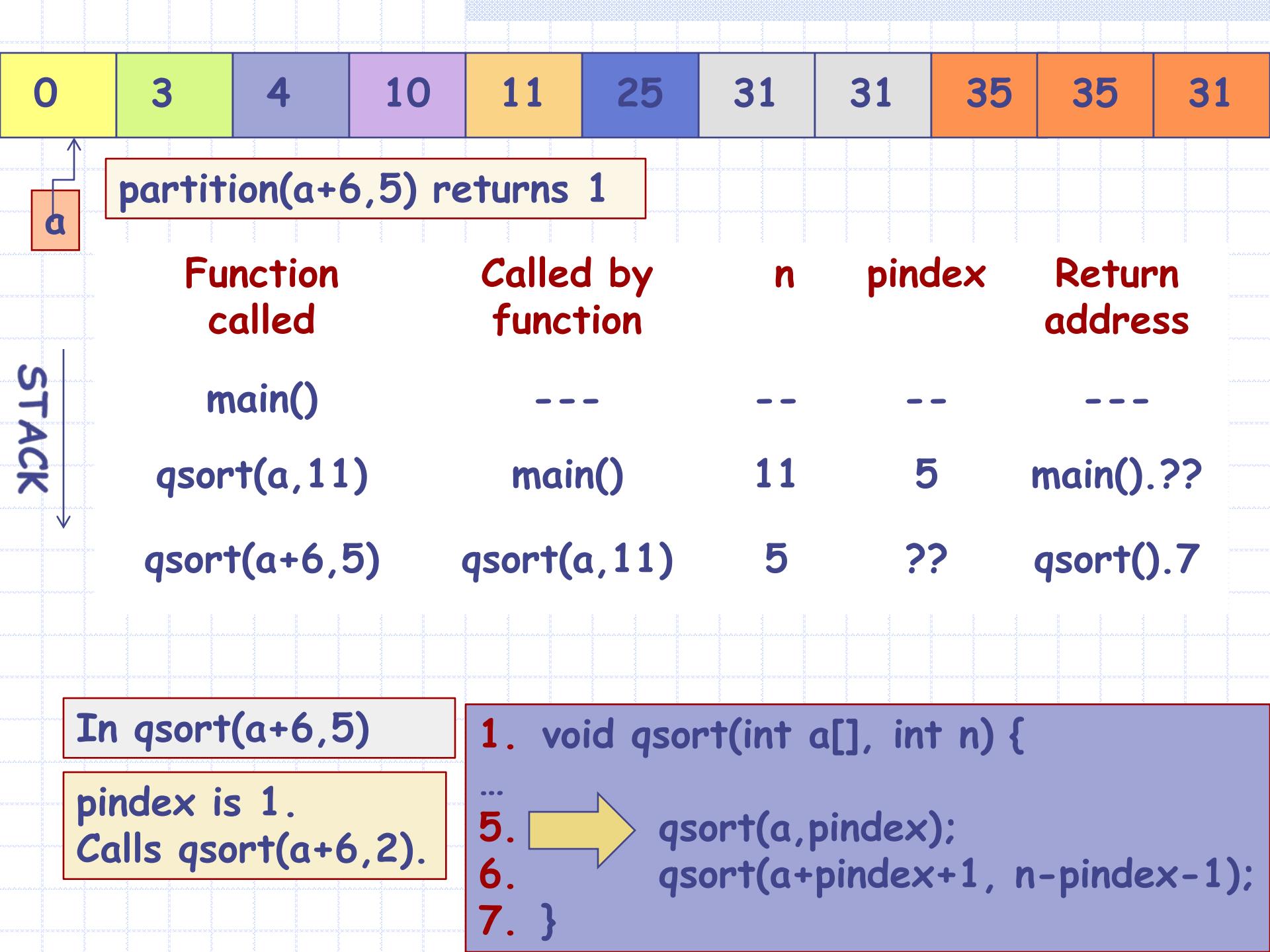
qsort(a, 6) resumes at line 7 and terminates. Control returns to calling fn qsort(a, 11) at line 6.

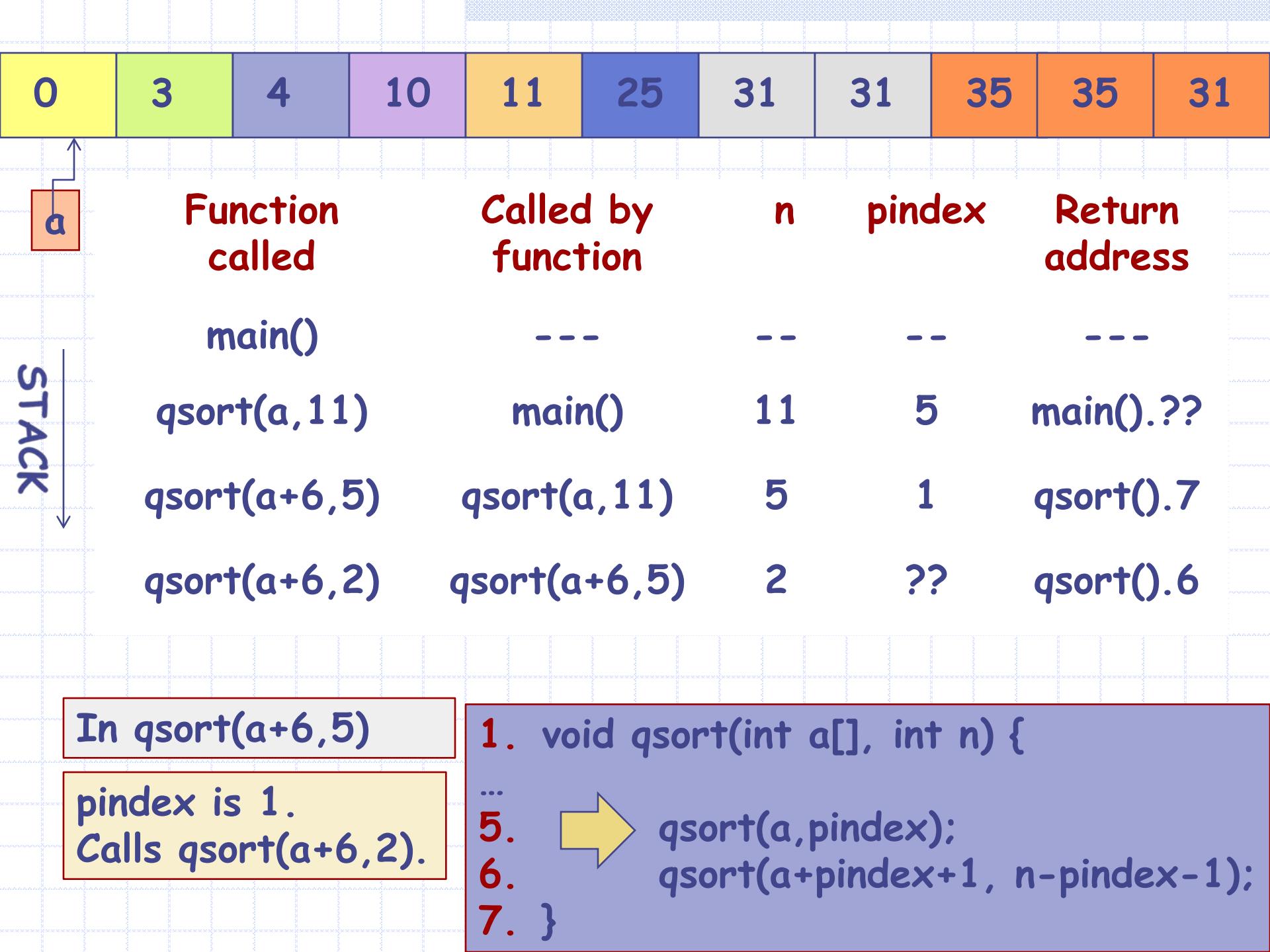
Note that qsort(a, 6) has terminated and the array a[0..5] has been sorted.

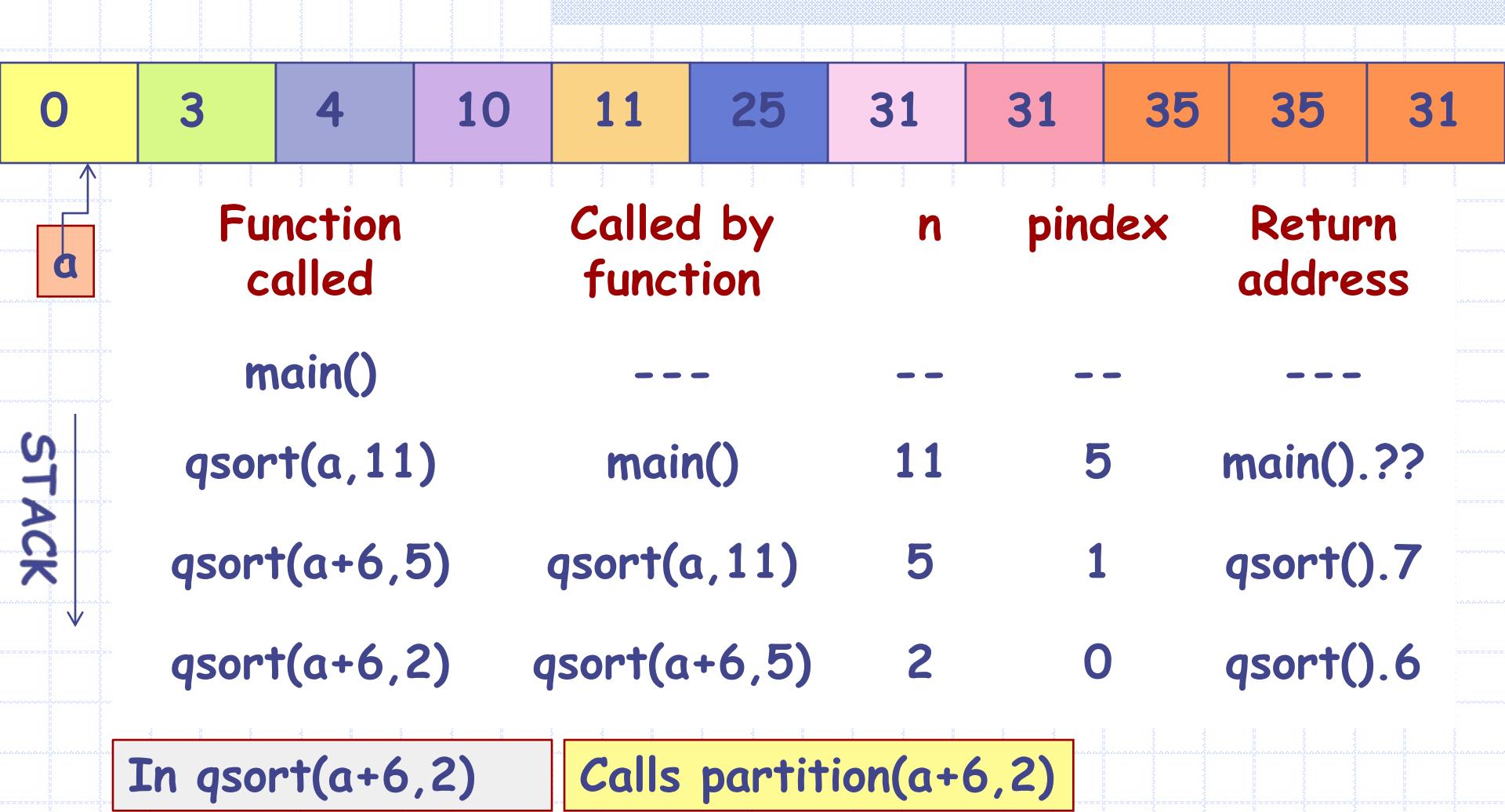


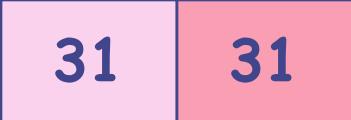


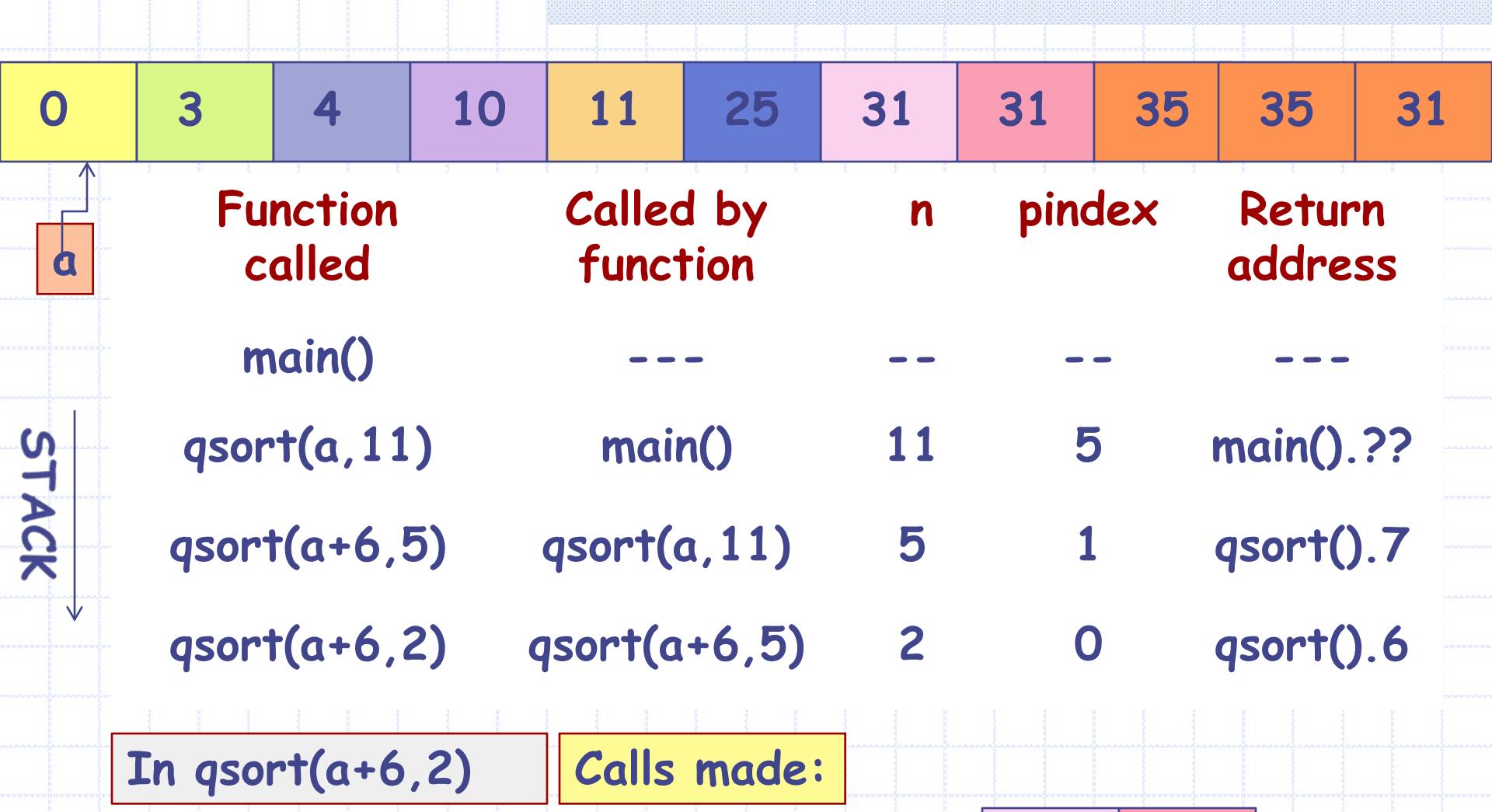




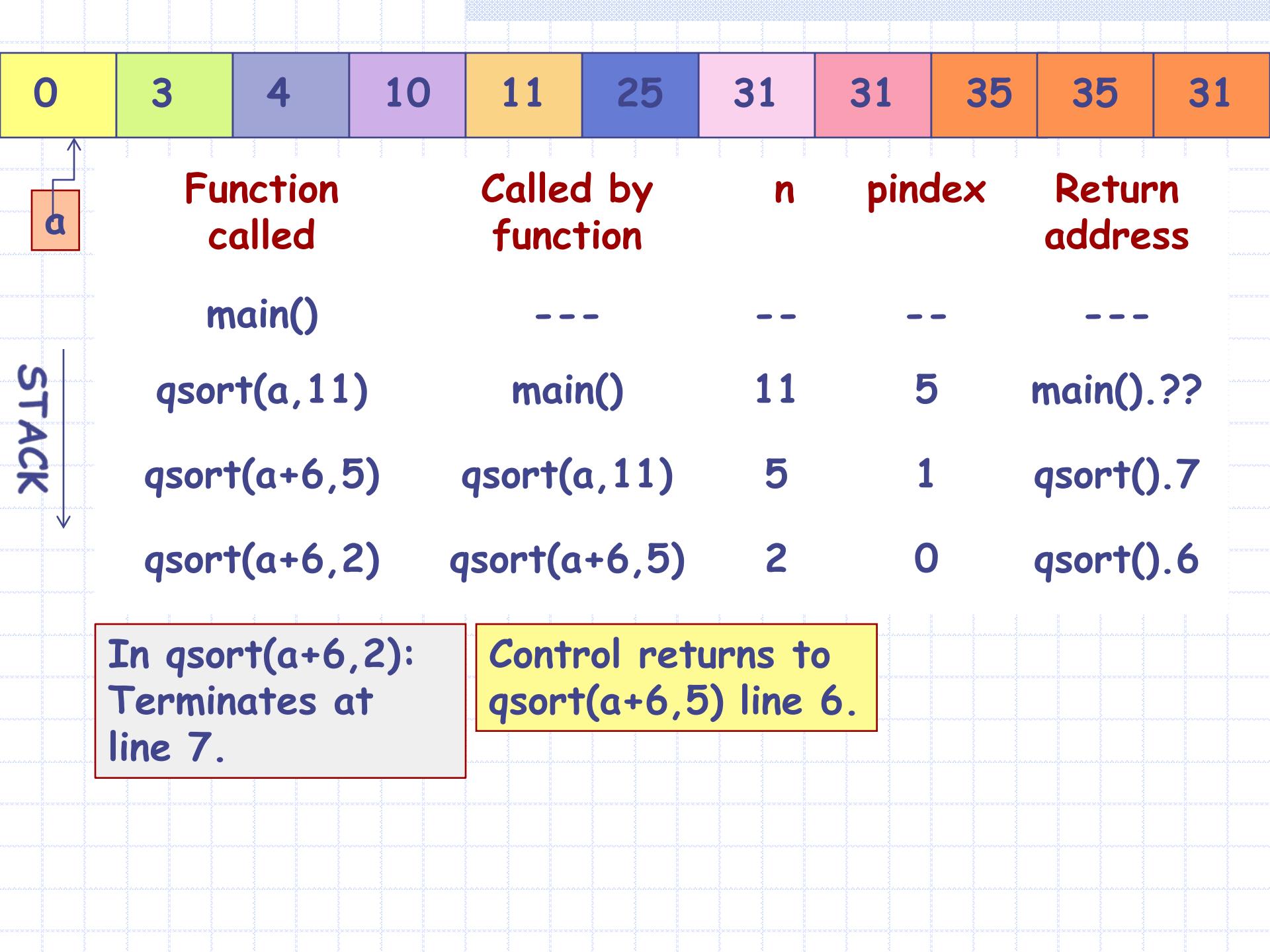




1. Partition is called for the array 
2. partition returns 0.
3. Control returns to qsort(a+6, 2) line 4, with pindex set to 0.
4. line 5: Calls qsort(a+6, 1) which returns immediately.
5. line 6: Calls qsort(a+7, 1) which returns immediately.



1. `partition(a+6, 2)` is called for the array
2. `partition(a+6, 2)` returns 0.
3. Control returns to `qsort(a+6, 2)` line 4, with `pindex` set to 0.
4. line 5: Calls `qsort(a+6, 1)` which returns immediately.
5. line 6: Calls `qsort(a+7, 1)` which returns immediately.



0	3	4	10	11	25	31	31	35	35	35	31
a											
Function called				Called by function		n	pindex	Return address			
main()				---		--	--	---			
qsort(a, 11)				main()		11	5	main().??			
qsort(a+6, 5)				qsort(a, 11)		5	1	qsort().7			

In qsort(a+6, 5) line 6.

calls qsort(a+8, 3)

```

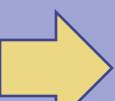
1. void qsort(int a[], int n) {
...
5.         qsort(a, pindex);
6.         qsort(a+pindex+1, n-pindex-1);
7.     }
    
```

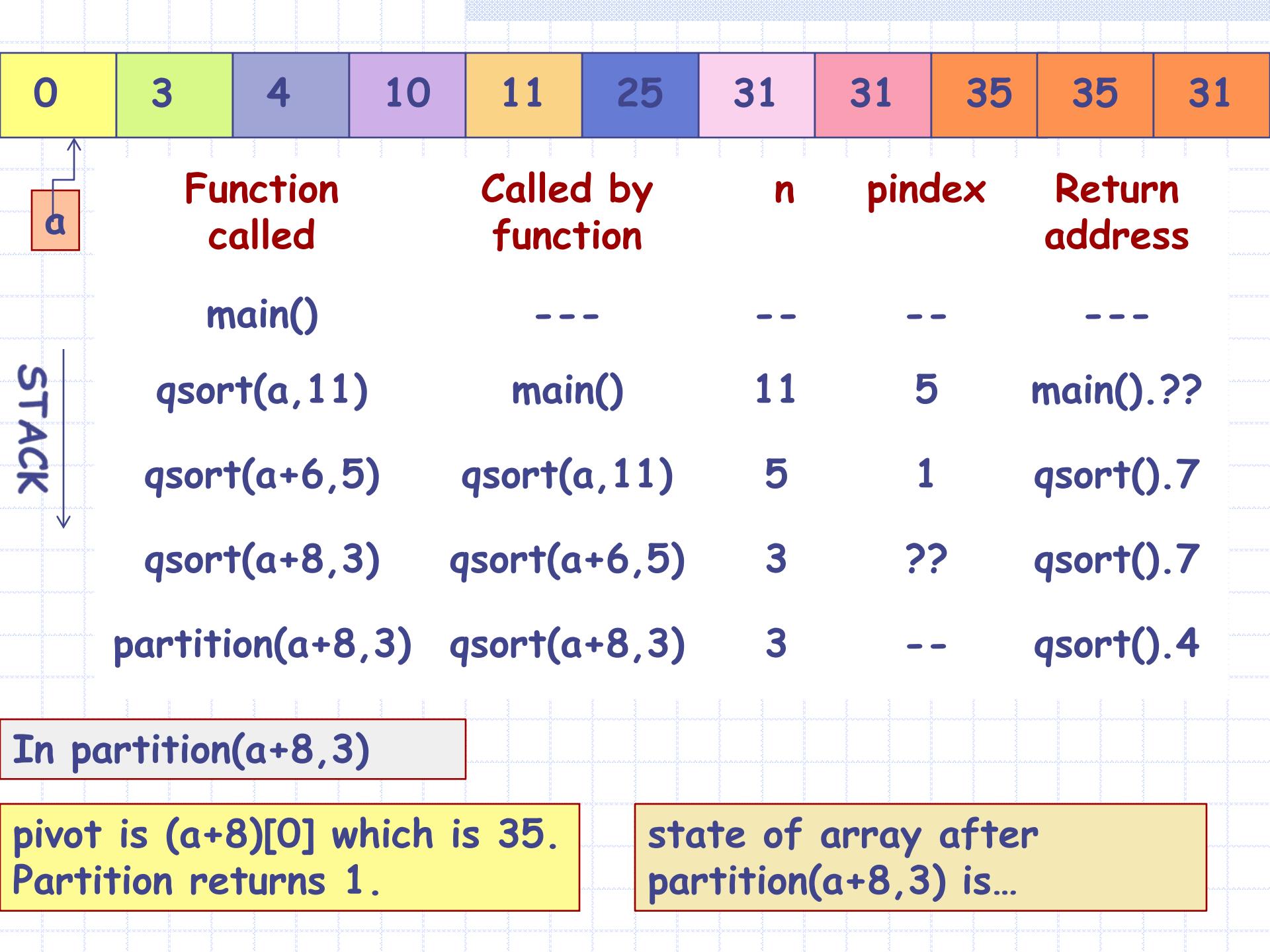
0	3	4	10	11	25	31	31	35	35	35	31
a											
Function called				Called by function		n	pindex	Return address			
main()				---		--	--	---			
qsort(a, 11)				main()		11	5	main().??			
qsort(a+6, 5)				qsort(a, 11)		5	1	qsort().7			
qsort(a+8, 3)				qsort(a+6, 5)		3	??	qsort().7			

In qsort(a+8, 3).

calls partition (a+8, 3)

1. void qsort(int a[], int n) {  
...  
5.     qsort(a, pindex);  
6.     qsort(a+pindex+1, n-pindex-1);  
7. }





0	3	4	10	11	25	31	31	31	35	35	
a											
Function called				Called by function		n	pindex	Return address			
main()				---		--	--	---			---
qsort(a, 11)				main()		11	5	main().??			
qsort(a+6, 5)				qsort(a, 11)		5	1	qsort().7			
qsort(a+8, 3)				qsort(a+6, 5)		3	1	qsort().7			
qsort(a+8, 2)				qsort(a+8, 3)		2	??	qsort().6			
partition(a+8, 2)				qsort(a+8, 2)		1	--	qsort().4			

In partition(a+8, 2)

1. pivot is 31.
2. returns 0. No change to array.

0	3	4	10	11	25	31	31	31	35	35
a										
Function called				Called by function		n	pindex		Return address	
main()				---		--	--		---	
qsort(a, 11)				main()		11	5	main().??		
qsort(a+6, 5)				qsort(a, 11)		5	1	qsort().7		
qsort(a+8, 3)				qsort(a+6, 5)		3	1	qsort().7		
qsort(a+8, 2)				qsort(a+8, 3)		2	??	qsort().6		

In qsort(a+8, 2) line 4:  
 1. calls partition(a+8, 2).

1. pivot is 31.
2. returns 0. No change to array.

0	3	4	10	11	25	31	31	31	35	35
a										
Function called				Called by function		n	pindex		Return address	
main()				---		--	--		---	
qsort(a, 11)				main()		11	5	main().??		
qsort(a+6, 5)				qsort(a, 11)		5	1	qsort().7		
qsort(a+8, 3)				qsort(a+6, 5)		3	1	qsort().7		
qsort(a+8, 2)				qsort(a+8, 3)		2	0	qsort().6		

In `qsort(a+8, 2)`

- line 4: pindex is set to 0.
- line 5: calls `qsort(a+8, 1)`.
- this returns immediately.

- line 6: calls `qsort(a+9, 1)`.
- Returns immediately.
- line 7: `qsort(a+8, 2)` returns.

0	3	4	10	11	25	31	31	31	35	35
a										
Function called				Called by function		n	pindex		Return address	
main()				---		--	--		---	
qsort(a, 11)				main()		11	5	main().??		
qsort(a+6, 5)				qsort(a, 11)		5	1	qsort().7		
qsort(a+8, 3)				qsort(a+6, 5)		3	1	qsort().7		
qsort(a+8, 2)				qsort(a+8, 3)		2	0	qsort().6		

STACK

In qsort(a+8, 2)

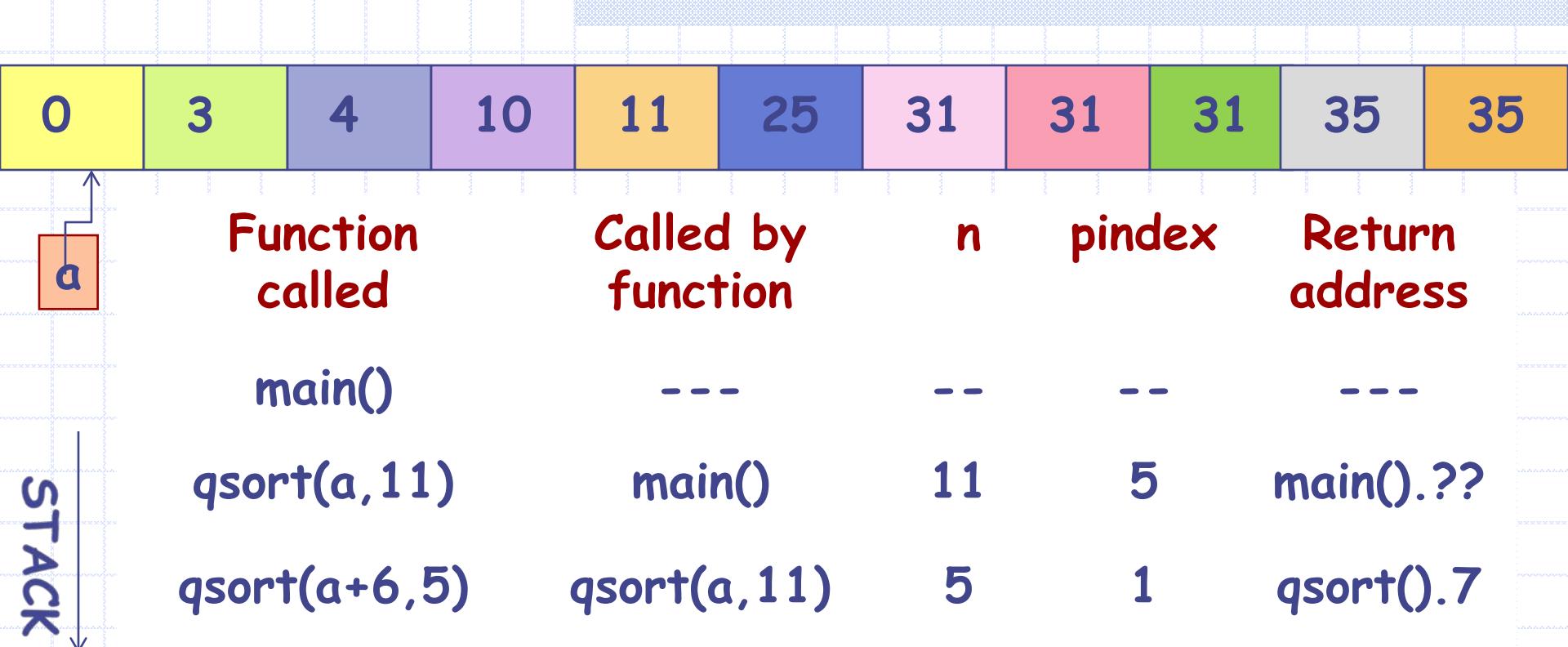
- line 4: pindex is set to 0.
- line 5: calls qsort(a+8, 1).
- this returns immediately.

- line 6: calls qsort(a+9, 1).
- Returns immediately.
- line 7: qsort(a+8, 2) returns to call qsort(a+8, 3) line 6.

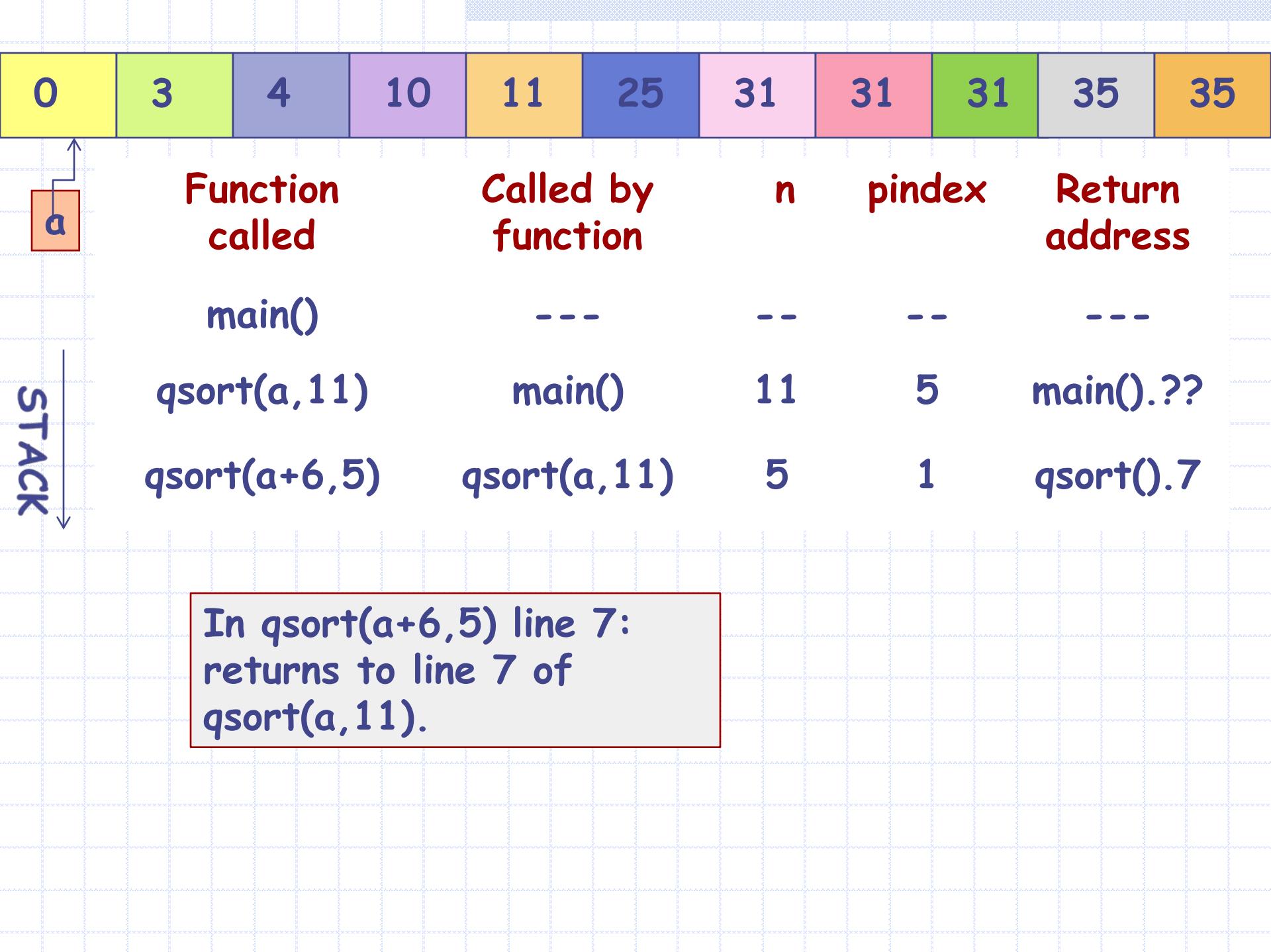
0	3	4	10	11	25	31	31	31	35	35
a										
Function called				Called by function		n	pindex		Return address	
main()				---		--	--		---	
qsort(a, 11)				main()		11	5	main().??		
qsort(a+6, 5)				qsort(a, 11)		5	1	qsort().7		
qsort(a+8, 3)				qsort(a+6, 5)		3	1	qsort().7		

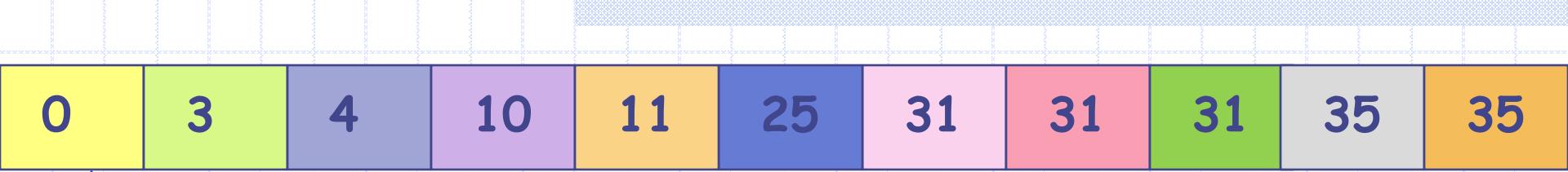
In qsort(a+8, 3) line 6:

1. calls qsort(a+10, 1)
2. returns immediately.
3. qsort(a+8, 3) returns to line 7 in call qsort(a+6, 5)



In `qsort(a+6, 5)` line 7:  
returns to line 7 of  
`qsort(a, 11)`.





a	Function called	Called by function	n	pindex	Return address
main()	---	--	--	--	---
qsort(a, 11)	main()	11	5	main().??	

In qsort(a, 11) line 7:  
returns to the calling  
function main().

